

RRRRRRRR	MM	MM	333333	GGGGGGGG	EEEEEEEEE	TTTTTTTT
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RR RR	RR	MMMM	33 33	GG	EE	TT
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1 0001 0 MODULE RM3GET (LANGUAGE (BLISS32) .
2 0002 0 IDENT = 'V04-000' .
3 0003 0 ) =
4 0004 1 BEGIN
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6 0006 1 ****
7 0007 1 *
8 0008 1 * COPYRIGHT (c) 1978, 1980, 1982, 1984 BY
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27 0027 1 ****
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29 0029 1 ++
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31 0031 1 FACILITY: RMS32 INDEX SEQUENTIAL FILE ORGANIZATION
32 0032 1
33 0033 1 Abstract: This module implements the get and find record operations
34 0034 1 for the indexed file organization.
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38 0038 1 ENVIRONMENT:
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40 0040 1 VAX/VMS OPERATING SYSTEM
41 0041 1
42 0042 1 --
43 0043 1
44 0044 1
45 0045 1 AUTHOR: E. H. MARISON CREATION DATE: 18-APR-78 13:11
46 0046 1
47 0047 1
48 0048 1 MODIFIED BY:
49 0049 1
50 0050 1 V03-025 Jim Teague 13-Aug-1984
51 0051 1 Fix bug in re-accessing records after they have been
52 0052 1 found to be locked. If a process had done a $GET on
53 0053 1 a record, then a $RELEASE, and then had to wait to
54 0054 1 $GET the record a second time, too much context was
55 0055 1 still around from the first $GET. This caused problems
56 0056 1 when the sought-after record had been deleted. RMS
57 0057 1 treated the $GET + $RELEASE + $GET case just like
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a \$FIND + \$GET case, and would end up with the wrong record.

Also, improve \$GET/\$FIND performance. Leave the infinite GET_RECORD loop immediately if GET_RECORD returns an unqualified success status. Formerly, RMS was forced to grind through an unbelievably perverted IF test EVERY TIME it returned from GET_RECORD.

V03-024 TSK0001 Tamar Krichevsky 15-Jun-1983
Change addressing mode to long relative for RM\$RU_RECLAIM.

V03-023 MCN0015 Maria del C. Nasr 24-Mar-1983
More linkages reorganization.

V03-022 TMK0015 Todd M. Katz 11-Mar-1983
If RMS had to wait for a record lock, and it must re-position to the primary data record by calling RMS\$FIND_BY_RRV, then make sure the primary data bucket containing the record is locked exclusively if the possibility exists that some reclamation maybe done (the file is write accessed and RU Journallable).

V03-021 MCN0014 Maria del C. Nasr 24-Feb-1983
Reorganize linkages

V03-020 TMK0014 Todd M. Katz 14-Jan-1983
Add support for Recovery Unit Journalling and RU ROLLBACK Recovery of ISAM files. Support involves modifications to RMS\$GET3B and RMS\$GET_RECORD.

The purpose of the routines within this module is to retrieve a non-deleted primary data record by the user specified access mode. If during its search for such a record RMS in its low-level routines encounters records that are marked RU_DELETE, RMS will try and delete them for good at this time provided it has write access to the file and the Recovery Unit in which they were deleted has completed successfully.

If RMS is able to delete a primary data record marked RU_DELETE in these low-level routines, then RMS proceeds to continue looking for a non-deleted primary data record just as if it had encountered a deleted record in the first place. Likewise, if RMS is unable to delete a record that is marked RU_DELETE because it does not have write access to the file, it merely continues its search. However, if RMS is unable to delete the record for good in these low-level routines because the Recovery Unit in which it was marked RU_DELETE has not successfully terminated, then RMS returns this record as if it was the non-deleted primary data record to be returned, and lets a higher-level routine decide whether or not to wait for the Recovery Unit in which the record was deleted to complete, or to return an error to the user.

The routines within this module are the high-level routines which decide what to do with RU_DELETED records that are returned from the low-level positioning routines.

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1. If RMS is unable to lock such a record because another process currently has it locked, then an RLK error is returned.
2. If RMS is able to lock such a record, regardless of whether it had to wait for it or not, then if it finds that the record is not marked RU_DELETE it will return it provided all other normal conditions have been met.
3. If on the other hand, RMS finds that the record is still marked RU_DELETE after it has locked it, then it will delete the record for good at this time (if the stream has write access to the file), and continue the search for a non-deleted primary data record provided the access mode is not by RFA.

I have also made two other changes in support of RU Journaling and Recovery. First, the ROP bit RABSV_NLK is totally ignored whenever a stream is currently within a Recovery Unit. Finally, it is also possible that a RU_UPDATE marked record might be re-formatted before releasing the bucket in which it is found provided the stream has write access to the file. The record being re-formatted in this case can only be the record that is to be returned as the non-deleted primary data record.

I have made an additional change to RMSGET RECORD. If RMS is currently randomly positioning by key to what it thinks is the current record, then it query locks the current record to make sure that this record is in fact locked to avoid a window in which the record is deleted between the time the record lock is released, and the bucket in which the record is found is accessed. If the user has specified record waiting it is disabled for this query lock. Currently it is disabled by clearing the RABSV_WAT bit if it is set, and then re-establishing its state after the query lock. The state bit IRBSV_NO_Q_WAIT maybe set to accomplish this same thing and it avoids modifying the user's control block.

I have created a routine RMSPOS_RFA whose functionality parallels that of RMSPOS_SEQ and RMSPOS_KEY. That is, the routine RMSGET_RECORD will call this routine whenever it is to position to the next primary data record by RFA instead of performing the positioning itself.

V03-019 TMK0013 Todd M. Katz 09-Nov-1982
Fix a bug in record unlocking. Whenever RMS must wait for a record lock (the RABSV_WAT ROP bit is set), and upon being granted the lock finds that the record it has been waited on has been deleted, RMS must perform a re-positioning. (There is one exception to this rule. If RMS was accessing the record by its RFA then the record deleted error is returned.) RMS must also perform a re-positioning whenever it is positioning by means of an alternate key and has had to wait for a record lock. As part of this re-positioning, RMS must release the lock it obtained during the prior positioning attempt. The problem is that RMS was using the wrong RFA when it went to release the record

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whenever it was re-positioning because the record it had to wait for had been deleted while it was waiting for it. The RFA it was using was the RFA of the current record. This record had been locked during the previous positioning operation, and had been unlocked during the first positioning attempt of the current operation. The lock RMS wants to release is for the record locked during the previous positioning attempt.

The fix for this problem is relatively straightforward. RMS never has to re-position unless it has had to wait for a record lock. Therefore, what I did was set the state bit OK_WAT_STATUS whenever a re-positioning has to be done instead of just setting it whenever the re-positioning is being done because RMS had to wait for a record lock while positioning along an alternate index. The setting of this state bit forces RMS to unlock the correct record during the re-positioning attempt.

V03-018 TMK0012 Todd M. Katz 29-Oct-1982
Make sure that RMS has the index descriptor for the primary key before the size of the primary data record to be returned is determined, the record unpacked (if the file is a prologue 3 file), and the record moved into the user's buffer.

V03-017 TMK0011 Todd M. Katz 11-Oct-1982
Fix a record locking bug. Whenever the ROP bit RAB\$V_WAT is set the possibility exists that RMS might have to wait for a record lock. If RMS is positioning by means of an alternate index, and has to wait for such a record lock, then it had to give up the SIDR bucket while it was waiting. Because it gave up the SIDR bucket, the information which it has in order to update the NRP list can no longer be considered valid. Since there is no way for RMS to easily re-access the SIDR bucket, RMS must re-position to it by re-calling GET_RECORD. Part of this re-positioning includes unlocking the very same primary data record which it had to wait for a record lock on. Unfortunately, GET_RECORD uses the NRP information to unlock primary data records, and RMS of course, didn't get to the point where it updated the NRP! Therefore, RMS is either not unlocking any record, or it is unlocking the wrong record. Both cases represent errors.

To fix this what I have done is added an input parameter to GET_RECORD. If it is set, RMS is re-positioning because of the above mentioned problem, and uses the RFA internally saved from the prior positioning attempt to unlock the record; otherwise, the RFA from the current record saved as part of the NRP context is used. Furthermore, whenever RMS does such a re-positioning, it now notes that it had to do so because of an OK_WAT success status positioning on an alternate key of reference. If it is successful at re-positioning, it sets the status to OK_WAT which represents the status that it would have returned if the re-positioning had not been necessary.

I have made two additional changes concerning when re-positioning is required. First, if RMS is positioning by key value, and after waiting for a record lock finds that the record it has been waited for has been deleted, then RMS will perform a

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264 0264 1
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re-positioning to the next record which matches the search key in keybuffer 2 according to the characteristics of the search. Formerly, RMS was just returning a record deleted error, but I believe the other approach has more merit. Second, whenever RMS is positioning by an alternate key of reference (sequentially or randomly by key value), and must wait for a record lock, then RMS must re-position to re-establish the NRP information for the SIDR. Formerly, this re-positioning was not done if RMS was performing a random \$FIND. However, since the stream which has the record locked can delete the SIDR array positioned to by the waiting stream without deleting the actual primary data record (by means of an \$UPDATE), then as the record eventually returned would not have the "correct" alternate key if re-positioning were not done, I believe that this requires this re-positioning to take place, even though the NRP is not going to be updated by this particular operation.

Finally, the last thing I did was make some changes on how the record unlocking is done when buffer errors are encountered during a \$GET/\$FIND. At this point the record has already been locked, and must be unlocked before control returns to the user. The routine GET_RECORD returns information in AP to RMS\$GET3B as to whether any special action is required to unlock this record on buffer errors. Unfortunately, AP is used throughout the remainder of RMS\$GET3B as input to record unpacking and key extraction; thus, its contents should a buffer error be detected and the record need to be unlocked, are unreliable. To fix this problem, I now set a flag bit on return from GET_RECORD if in fact special action will be required to unlock the record on buffer errors, and reference this bit in that circumstance rather than the AP.

V03-016 TMK0010 Todd M. Katz 29-Sep-1982
If a file is a prologue 3 file with alternate keys, and RMS is positioning by means of an alternate key of reference, then RMS was not unpacking the record before returning it to the user because it assumed that the record had been unpacked during the positioning and there was no need to unpack it a second time. However, while this is true, RMS at this point does not know the unpacked record's size. Thus, for the time being RMS must always unpack the record before moving it into the user's record buffer if the file is a prologue 3 file.

V03-015 TMK0009 Todd M. Katz 09-Sep-1982
The field IRB\$B_SRCHFLAGS is now a word in size. Change all references to it.

Whenever RMS is positioning by means of an alternate key of reference (IRB\$B_RP_KREF > 0), then there is never a need in the local routine GET_RECORD to extract the alternate key of the record positioned to into keybuffer 2. This is because as part of positioning to the primary data record from the SIDR in the first place, the SIDR key has already been extracted into keybuffer 2.

Eliminate all references to the routine RMSKEY_TYPE_CONV, since this routine doesn't do anything anyway.

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340 0340 1
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The only time it is necessary to check for a valid packed decimal key is when the key type is packed decimal. It is never necessary to check for a valid packed decimal type when there is more than one segment and the file is a prologue 3 file. The packed decimal verification routine no longer requires parameters.

V03-014 KBT0294 Keith B. Thompson 23-Aug-1982
Reorganize psects

V03-013 TMK0008 Todd M. Katz 10-Aug-1982
At the present time, when the accessing of a record by RFA fails, the error returned by RMS\$IND_BY_RRV is the error that gets reported to the user. Change this so that if this routine returns an error of RMSS_EOF (because the RFA VBN is greater than the VBN of any primary data bucket), this error gets mapped into an error of RMSS_PNF.

V03-012 MCN0013 Maria del C. Nasr 10-Aug-1982
Check for less than 0 on call to RMS\$COMPARE_KEY so that LIM check is done correctly. This is to fix bug introduced by MCN0012.

V03-011 TMK0007 Todd M. Katz 19-Jun-1982
Implement the RMS cluster NRP solution. Basically this involves removal of the NRP cells from system space, and the maintenance of the next record positioning context locally within the IRAB. Changes required to the routines in this module are as follows:

1. The routine SETUP_NRP_DATA now sets up the current record context in the process local IRAB instead of in the system-wide NRP cell.
2. The IRAB variables IRBSL_NEXT_VBN and IRBSW_NEXT_ID are used to temporarily hold the RFA address of the "next" primary data record until the updating of the local NRP context takes place. This is because nothing in the local NRP context maybe modified, until everything is modified!
3. The local routines must also be modified both to make use of the next record positioning context now saved within the IRAB instead of within a systemwide NRP cell.
4. If RMS encounters the end-of-file set the IRBSV_EOF bit. This bit is also cleared after successfully positioning randomly by key value. The former function of this bit has now been taken over by the new bit IRBSV_CON_EOF.
5. Special processing is required for \$GETs following random \$FINDs. A random \$FIND does not change the notion of what the next record is although it does change the notion of what the current record is! Example with the record sequence 0 A B - sequential \$GET to A, random \$FIND to 0, \$DELETE 0, followed by a sequential \$GET returns B, the next record. The random \$FIND changed the current record to 0, but did not change the next record to 0! The RMS cluster solution

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399 0399 1

for NRP positioning handles this by keeping the current primary data record's RFA and the RFA of the primary data record for NRP positioning in separate fields. Most operations set all NRP fields and as a result the RFA address of the current primary data record and the RFA address of the primary data record used for NRP positioning are the same. However, a random \$FIND will set only the current primary data record's RFA field. If the random \$FIND is immediately followed by a sequential \$GET, then it is only at that moment that the local NRP context is setup to return the randomly found record as the next record.

Also, it is no longer necessary within GET RECORD to loop on calls to RM\$POS_SEQ or RM\$POS_KEY when these routines return RLK errors. An RLK error could occur only when positioning on an alternate index and signalled that re-positioning should be forced. This re-positioning is now handled at a much lower level, and there is no longer any need to force it.

During the performance optimization of TMK0005 one incorrect assumption was made: that no deleted records were encountered between the last record retrieved, whose key is in keybuffer 1, and the new record that has just been retrieved. If this is true, the optimization holds, but if it is not, we can not use the key of the last retrieved record to uncompress the key of the new record, because the compression of the key of the new record is based upon the intervening deleted records, and not the key of the last record. In such a situation, the key of the new record must be extracted, and re-expanded in the old way performing a bucket scan if necessary.

Finally, it will no longer be necessary to unpack the primary data record when the file is a prologue 3 file, and RMS is currently positioning by an alternate key since the record will have been already unpacked and is within the internal record buffer.

V03-010 MCN0012 Maria del C. Nasr 29-Jun-1982
Allow keys of different data types other than string.
Change all CH\$COMPARE calls to RM\$COMPARE_KEY to compare
keys taking into consideration the different data types.

V03-009 TMK0006 Todd M. Katz 26-May-1982
I have changed how the ROP=LIM key comparison is performed. Formerly, the routine RM\$COMPARE_REC was being called. It was being called because the (incorrect) assumption existed that the key of the next record might have to be extracted and re-expanded, if key compression was enabled, in order to make the comparison. As it turns out, at this point in the operation, RMS has already extracted (and re-expanded if necessary) the key of the next record into keybuffer 2. Thus, in order to make the comparison, only a CH\$COMPARE between the key in the user's key buffer and the key in keybuffer 2 need be made. Thus, this comparison has now been made prologue independent, it is a performance optimization for all prologue versions, and the performance realized for prologue 3 files is considerable because it eliminates the need for one more

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bucket scan which was unnecessarily being done to re-expand the key of the next record.

V03-008 TMK0005 Todd M. Katz 26-May-1982
Performance enhancement. After successfully positioning to the next record, RMS extracts its key into keybuffer 2. If key compression is enabled this mandates another bucket pass to re-expand the key. However, if RMS is positioning sequentially, then it has the key of the previous record retrieved saved in keybuffer 1. RMS can use this key to supply any characters front compressed off the key of the current record instead of performing another bucket pass to expand the key.

V03-007 TMK0004 Todd M. Katz 24-May-1982
Performance enhancement. When performing a \$GET on a prologue 3 file, the record found must be unpacked before it is returned. Part of this unpacking includes extraction of the primary key from its position in front of the data record, and its re-expansion if key compression is enabled. But if we are positioning by primary key of reference then there is really no need to extract and re-expand the primary key because RMS already has it in the proper form within keybuffer 2. To signal to the routine RMSUNPACK_REC, that there is no need to extract and re-expand the primary key of the found data record, but that it maybe found in keybuffer 2, we initialize AP to 2 before calling the routine when the key of reference is the primary key.

V03-006 LJA0008 Laurie Anderson 08-Apr-1982
Must check for allocation of IDX_DFN, before access fields in it. The IDX_DFN will not be allocated if there is an error returned from RMSKEY_DESC which is called by GET_RECORD.

V03-005 TMK0003 Todd M. Katz 01-Apr-1982
If record locking is enabled, in GET_RECORD we lock the record we have found. If we had to wait for this record, the status returned is an alternate success (OK_WAT). We should be setting the IRBSV_UNLOCK_RP bit so that whenever we have finished with this record RMS will know to release it, but because our status is OK_WAT and not success, the current flow of control forces a return before this bit can be set. Therefore, the possibility exists that once a process has waited for a record lock and successfully locks the record, it will not release the lock unless explicitly told to do so (such as by an explicit \$RELEASE). To avoid this undesirable possibility, we will make sure that IRBSV_UNLOCK_RP will be set even when we had to wait for a record lock.

V03-004 TMK0002 Todd M. Katz 26-Mar-1982
Under two different set of circumstances we will have to release the record lock obtained in GET_RECORD.

1. If we have decided to make another iterative call to GET_RECORD and we have locked a record within the last call, then we must release this lock before attempting to locate the next record in the current call.

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461 0461 1
462 0462 1
463 0463 1
464 0464 1
465 0465 1
466 0466 1
467 0467 1
468 0468 1
469 0469 1
470 0470 1
471 0471 1
472 0472 1
473 0473 1
474 0474 1
475 0475 1
476 0476 1
477 0477 1
478 0478 1
479 0479 1
480 0480 1
481 0481 1
482 0482 1
483 0483 1
484 0484 1
485 0485 1
486 0486 1
487 0487 1
488 0488 1
489 0489 1
490 0490 1
491 0491 1
492 0492 1
493 0493 1
494 0494 1
495 0495 1
496 0496 1
497 0497 1
498 0498 1
499 0499 1
500 0500 1
501 0501 1
502 0502 1
503 0503 1
504 0504 1
505 0505 1
506 0506 1
507 0507 1
508 0508 1
509 0509 1
510 0510 1
511 0511 1
512 0512 1
513 0513 1

2. If we are currently performing a random SFIND/SGET, and we must wait in our attempt to lock the record we have found (RABSV_WAT is set), and upon returning and reaccessing the bucket we found that this record has been deleted by the stream that previously had it locked, then we must release our lock on this deleted record before returning the status of deleted record from RM\$GET3B.

Both of these record lock releases may be signaled by setting the IRAB bit IRBSL_UNLOCK_RP within GET_RECORD at the appropriate time. The record will then be locked either within GET_RECORD in the former case or within RM\$GET3B in the latter.

V03-003 TMK0001 Todd M. Katz 24-Mar-1982
Change all references to te keybuffers to use the macro KEYBUF_ADDR

If an error status of RLK is returned on an attempt to "get" a data record, try again until the record is retrieved or a different error is returned. This will only occur when our key of reference is other than key 0, and someone else had the primary data (or RRV) bucket locked when we attempted to access it from the SIDR. The SIDR bucket must be released and we have to reaccess it inorder to avoid a potential deadlock situation, and returning an error of RLK will now guarantee that this is what will happen.

If the attempt to sequentially access a record results in a status of record deleted being returned from GET_RECORD, attempt to sequentially retrieve the very next record, and continue doing this until some other status is returned. This situation can develop if we try for a record lock and end up waiting (the ROP_WAT bit is set) for it. While we are waiting the process (or stream) which has the record lock deletes it. When control returns to this process, the status it gets back indicates that it had to wait, and so it reaccesses the bucket the record was in (it had to release it when it went for the record lock) and now finds the record is deleted and returns that status.

When control returns to RM\$GETRECORD from RM\$POS_SEQ, RM\$POS_KEY, or RM\$FIND_BY_RRV with a success status, the next record has been found, the bucket containing it has been locked, and the IRAB fields IRBSL_RFA_VBN, IRBSW_RFA_ID, and IRBSW_SAVE_POS contain the information necessary to update the NRP context to that of the "found" record. If a decision is made to lock the record, and RMS has to stall for the record lock (RABSV_WAT is set and some other stream has the record locked) then when the lock is obtained, the bucket containing the record is no longer locked (if we have to wait for the record lock we must release the lock on the bucket to avoid the possibility of deadlock), and the NRP updating information within the IRAB can no longer be considered valid because the bucket containing the record might have split moving the record to the new bucket, and the record itself might even have been deleted. If we are accessing this file by its primary key, then as its record pointer (RP) information is still valid,

514 0514 1 |
515 0515 1 |
516 0516 1 |
517 0517 1 |
518 0518 1 |
519 0519 1 |
520 0520 1 |
521 0521 1 |
522 0522 1 |
523 0523 1 |
524 0524 1 |
525 0525 1 |
526 0526 1 |
527 0527 1 |
528 0528 1 |
529 0529 1 |
530 0530 1 |
531 0531 1 |
532 0532 1 |
533 0533 1 |
534 0534 1 |
535 0535 1 |
536 0536 1 |
537 0537 1 |
538 0538 1 |
539 0539 1 |
540 0540 1 |
541 0541 1 |
542 0542 1 |
543 0543 1 |
544 0544 1 |
545 0545 1 |
546 0546 1 |
547 0547 1 |
548 0548 1 |
549 0549 1 |
550 0550 1 |
551 0551 1 |
552 0552 1 |
553 0553 1 |
554 0554 1 |
555 0555 1 |
556 0556 1 |
557 0557 1 |
558 0558 1 |
559 0559 1 |
560 0560 1 |
561 0561 1 |
562 0562 1 |
563 0563 1 |
564 0564 1 |
565 0565 1 |
566 0566 1 |
567 0567 1 |
568 0568 1 |
569 0569 1 |
570 0570 1 |

we can call RMS\$IND_BY_RRV to lock the bucket. Once the bucket has been again locked, the information necessary to update the next record context can be obtained and we can proceed. However, if we are accessing this record by an alternate key and we stall, then there is no easy way to reobtain the next record context information necessary for later updating of the NRP list. Thus, if the primary data record which we have locked is not deleted, and if we are performing a sequential \$FIND, or a \$GET operation (for which NRP list updating is mandatory), then we have no choice but repeat this lengthy process from the beginning. This is done BY by noting on return from GET_RECORD, that we have an alternate success status (OK_WAT), that we our key of reference is not the primary key, and that we have not locked the primary data bucket.

V03-002 LJA0006 Laurie Anderson 23-Mar-1982
If the get record caused an RTB error, then the primary key was not copied into an RMS internal buffer. This key buffer is used to avoid un-locking a record during a random access for an exact match by key when that record is the current record.

V03-001 KPL0009 Peter Lieberwirth 17-Mar-1982
Set UNLOCK_RP on errors reaccessing record after successful wait for record lock. This will cause the record to be unlocked on the way out.

Add subtitles.

V02-025 DJD0001 Darrell Duffy 1-March-1982
Clean up probing of input parameters

V02-024 KPL0008 Peter Lieberwirth 5-Nov-1981
Add support for PUT to EOF by clearing internal EOF flag on random gets, and returning RMSS_EOF on sequential gets.

V02-023 KPL0007 Peter Lieberwirth 7-Oct-1981
Fix bug on reaccessing buffer logic on secondary key.

V02-022 KPL0006 Peter Lieberwirth 2-Oct-1981
Fix bugs related to interaction of WAT bit set when QUERY_LCK is called. When re-accessing same record, don't WAT even if user said to until real lock logic. Also, when QUERY_LCK called with WAT set in lock logic of GET_RECORD, remember to reaccess the bucket if RMS stalled. Fix reaccess logic to work with secondary keys. (Oops!)

V02-021 KPL0005 Peter Lieberwirth 23-Aug-1981
Fix incorrect and misleading due to V02 v18.
Also, allocate a temporary variable more efficiently.

V02-020 MCN0011 Maria del C. Nasr 24-Jul-1981
Implement key type conversion.

V02-019 MCN0010 Maria del C. Nasr 23-Jul-1981
Incorporate all the following changes:

571 0571 1 Use RMSREC OVHD and user's buffer to get key
572 0572 1 Use key buffer 4 to unpack primary key.
573 0573 1 Include code for unpacking of prologue 3 data records.
574 0574 1 Change calling sequence of RMSFIND_BY_RRV.
575 0575 1 Increase size of record identifier to a word in the IRB,
576 0576 1 NRP, and RLB.
577 0577 1 Modify routine to handle new prologue 3 data structure
578 0578 1 changes (base level 1).
579 0579 1
580 0580 1 V02-018 KPL0004 P. Lieberwirth 15-Jan-1981 3:15
581 0581 1 Change GET RECORD to reaccess bucket if it had to be given
582 0582 1 up for wait on record. Implements new RDP functionality
583 0583 1 implied by WAT and RFA.
584 0584 1
585 0585 1 V02-017 SPR33597 P. Lieberwirth 24-Nov-1980 10:00
586 0586 1 Fix bug where omitted fetch operator caused incorrect test
587 0587 1 for validity of NRP. Bug caused incorrect operation on
588 0588 1 sequential \$FINDs. Clarify some commentary by cleaning up
589 0589 1 some spelling mistakes, and explaining FIND some more.
590 0590 1
591 0591 1 V02-016 REFORMAT K. E. Kinnear 24-Jul-1980 9:54
592 0592 1
593 0593 1 V02-015 CDS0073 C. D. Saether 17-Jan-1980 2:35
594 0594 1 Restructure current record unlocking logic to add check
595 0595 1 when duplicates aren't allowed to avoid record lock window
596 0596 1 and reaccessing current record.
597 0597 1
598 0598 1 V02-014 PSK0005 P. S. Knibbe 18-Dec-1979 5:00
599 0599 1 Check that packed decimal keys are in the correct format.
600 0600 1
601 0601 1 REVISION HISTORY:
602 0602 1
603 0603 1 V02-013 C. D. Saether 12-Jul-1979 11:30
604 0604 1 Level calling RMS\$COMPARE_REC should be -1.
605 0605 1
606 0606 1 V01-012 W. Koenig 6-Feb-1979 17:19
607 0607 1 Fill in user's RFA after some other checks.
608 0608 1
609 0609 1 V01-011 W. Koenig 6-Dec-1978 10:19
610 0610 1 Implement RMSS_OK_LIM.
611 0611 1
612 0612 1 V01-010 W. Koenig 5-Dec-1978 10:25
613 0613 1 Don't return DCT Field.
614 0614 1
615 0615 1 V01-009 W. Koenig 24-Oct-1978 14:02
616 0616 1 Make changes caused by sharing conventions.
617 0617 1
618 0618 1 V01-008 W. Koenig 5-Oct-1978 14:02
619 0619 1 Zero all the NRP flags when resetting the NRP data.
620 0620 1
621 0621 1 V01-007 W. Koenig 26-Sep-1978 16:42
622 0622 1 Don't zero the RP information after a successful get or
623 0623 1 sequential find.
624 0624 1
625 0625 1 V01-006 W. Koenig 26-Sep-1978 13:15
626 0626 1 Can no longer zero out RP_SECOND as a longword.
627 0627 1

628 0628 1 V01-005 C. D. Saether 21-Sep-1978 16:44
629 0629 1 Clear SRCFLAGS always.
630 0630 1
631 0631 1 V01-004 W. Koenig 21-Sep-1978 15:50
632 0632 1 Return the data to the user on any success, not just "suc".
633 0633 1
634 0634 1 V01-003 C. D. Saether 20-Sep-1978 16:25
635 0635 1 Clear NRP update flags when storing NRP.
636 0636 1
637 0637 1 V01-002 C. D. Saether 12-Sep-1978 15:21
638 0638 1 Remove NXTBDB setup on RFA access.
639 0639 1
640 0640 1 !*****
641 0641 1
642 0642 1 LIBRARY 'RMSLIB:RMS';
643 0643 1
644 0644 1 REQUIRE 'RMSSRC:RMSIDXDEF';
645 0709 1
646 0710 1 ! Define default psects for code
647 0711 1 !
648 0712 1 PSECT
649 0713 1 CODE = RM\$RMS3(PSECT_ATTR),
650 0714 1 PLIT = RM\$RMS3(PSECT_ATTR);
651 0715 1
652 0716 1 ! Linkages.
653 0717 1
654 0718 1 LINKAGE
655 0719 1 L_COMPARE_KEY,
656 0720 1 L_JSB,
657 0721 1 L_JSB01,
658 0722 1 L_PRESERVE1,
659 0723 1 L_QUERY_AND_LOCK,
660 0724 1 L_RABREG,
661 0725 1 L_RABREG_67,
662 0726 1 L_RABREG_7,
663 0727 1 L_REC_OVRD,
664 0728 1
665 0729 1 ! Local Linkages.
666 0730 1
667 0731 1 L_GET_RECORD = JSB () :
668 0732 1 GLOBAL (COMMON_RABREG, R_REC_ADDR, R_IDX_DFN)
669 0733 1 NOPRESERVE (2, 3, 4, 5),
670 0734 1 L_SETUP_NRP = JSB () :
671 0735 1 GLOBAL (COMMON_RABREG, R_IDX_DFN)
672 0736 1 NOPRESERVE (2, 3, 4, 5);
673 0737 1
674 0738 1 ! Forward Routines.
675 0739 1
676 0740 1 FORWARD ROUTINE
677 0741 1 GET_RECORD : L_GET_RECORD;
678 0742 1
679 0743 1 ! External Routines.
680 0744 1
681 0745 1 EXTERNAL ROUTINE
682 0746 1 RMSCOMPARE_KEY : RL\$COMPARE_KEY,
683 0747 1 RMSFIND_BY_RRV : RL\$RABREG_67,
684 0748 1 RMSKEY_DESC : RL\$RABREG_7,

685	0749	1	RMSLOCK	: RL\$QUERY_AND_LOCK,
686	0750	1	RMSNOREAD LONG	: RLSJSB,
687	0751	1	RMSNOWRT LONG	: RLSJSB,
688	0752	1	RMSPCKDEC CHECK	: RLSRABREG 7,
689	0753	1	RMSPOS KEY	: RLSRABREG 67,
690	0754	1	RMSPOS-RFA	: RLSRABREG 67,
691	0755	1	RMSPOS-SEQ	: RLSRABREG 67,
692	0756	1	RMSQUERY LCK	: RLSQUERY AND_LOCK,
693	0757	1	RMSRECORD_IC	: RLSRABREG 67,
694	0758	1	RMSRECORD_KEY	: RLSPRESERVE1,
695	0759	1	RMSRECORD_VBN	: RLSPRESERVE1,
696	0760	1	RMSREC OVHD	: RLSREC OVHD,
697	0761	1	RMSRLSBKT	: RLSPRESERVE1,
698	0762	1	RMSRU RECLAIM	: RLSRABREG 67 ADDRESSING_MODE(LONG_RELATIVE),
699	0763	1	RMSUNLOCK	: RL\$QUERY AND_LOCK,
700	0764	1	RMSUNPACK_REC	: RLSJSB01;
701	0765	1		

SETUP_NRP_DATA

703 0766 1 %SBTTL 'SETUP_NRP_DATA'
704 0767 1 ROUTINE SETUP_NRP_DATA : L_SETUP_NRP NOVALUE =
705 0768 1 !++
706 0769 1 !++
707 0770 1
708 0771 1 FUNCTIONAL DESCRIPTION:
709 0772 1
710 0773 1 This routine saves the next record positioning data
711 0774 1 in the IRAB from the temporary IRAB locations filled
712 0775 1 in during the positioning to the primary data record.
713 0776 1
714 0777 1 CALLING SEQUENCE:
715 0778 1
716 0779 1 SETUP_NRP_DATA()
717 0780 1
718 0781 1 INPUT PARAMETERS:
719 0782 1 NONE
720 0783 1
721 0784 1 IMPLICIT INPUTS:
722 0785 1
723 0786 1
724 0787 1 IRAB
725 0788 1 IRBSL_FIRST_ID - Current SIDR's first SIDR array element ID
726 0789 1 IRBSL_FIRST_VBN - Current SIDR's first SIDR array element VBN
727 0790 1 IRBSL_KEYBUF - Pointer to keybuffers (to access keybuffer 2)
728 0791 1 IRBSL_NEXT_ID - ID of current primary data record
729 0792 1 IRBSL_NEXT_VBN - VBN of current primary data record
730 0793 1 IRBSL_RFA_ID - ID of current record (SIDR/primary)
731 0794 1 IRBSL_RFA_VBN - VBN of current record (SIDR/primary)
732 0795 1 IRBSB_RP_KREF - Key of reference used to retrieve user data record
733 0796 1 IRBSW_SAVE_POS - Number of elements before current SIDR element
734 0797 1 IFBSW_KBUFSZ - Size of keybuffer (to access keybuffer 2)
735 0798 1
736 0799 1 OUTPUT PARAMETERS:
737 0800 1 NONE
738 0801 1
739 0802 1 IMPLICIT OUTPUTS:
740 0803 1
741 0804 1 IRAB
742 0805 1 IRBSW_CUR_COUNT - Number of elements before current SIDR element
743 0806 1 IRBSW_CUR_ID - ID of current record (SIDR/primary)
744 0807 1 IRBSB_CUR_KREF - Key of reference of current record (SIDR/primary)
745 0808 1 IRBSL_CUR_VBN - VBN of current record (SIDR/primary)
746 0809 1 IRBSV_EOF_ - clear indicating stream is not at end-of-file
747 0810 1 IRBSL_KEYBUF - Pointer to keybuffers (to access keybuffer 1)
748 0811 1 IRBSW_POS_ID - ID of primary data record for NRP positioning
749 0812 1 IRBSL_POS_VBN - VBN of primary data record for NRP positioning
750 0813 1 IRBSL_SIDR_VBN - Current SIDR's first SIDR array element VBN
751 0814 1 IRBSW_SIDR_ID - Current SIDR's first SIDR array element ID
752 0815 1 IRBSW_UDR_ID - ID of current primary data record
753 0816 1 IRBSL_UDR_VBN - VBN of current primary data record
754 0817 1 ROUTINE VALUE:
755 0818 1 NONE
756 0819 1
757 0820 1 --
758 0821 1
759 0822 2 BEGIN

```
SETUP_NRP_DATA

760 0823 2
761 0824 2 EXTERNAL REGISTER
762 0825 2 R IDX DFN STR,
763 0826 2 COMMON_RAB_STR;
764 0827 2
765 0828 2 ! Indicate that this stream is no longer at the file's end of file.
766 0829 2
767 0830 2 !RAB[IRBSV_EOF] = 0;
768 0831 2
769 0832 2 ! Move the VBN of the current record into the appropriate IRAB location
770 0833 2
771 0834 2 !RAB[IRBSL_CUR_VBN] = .IRAB[IRBSL_RFA_VBN];
772 0835 2
773 0836 2 ! If the current record happens to also be the primary data record, then
774 0837 2 move its ID into the appropriate IRAB location.
775 0838 2
776 0839 2 IF .IRAB[IRBSB_RP_KREF] EQLU 0
777 0840 2 THEN
778 0841 2 !RAB[IRBSW_CUR_ID] = .IRAB[IRBSW_RFA_ID]
779 0842 2
780 0843 2 ! If the current record happens to be a SIDR, then it has no ID to save,
781 0844 2 and instead save the SIDR first array element's VBN and ID (this uniquely
782 0845 2 identifies the SIDR), and the number of array elements preceding the
783 0846 2 current element (which points to the primary data record that is being
784 0847 2 retrieved).
785 0848 2
786 0849 2
787 0850 3 ELSE
788 0851 3 BEGIN
789 0852 3 !RAB[IRBSW_CUR_COUNT] = .IRAB[IRBSW_SAVE_POS];
790 0853 3 !RAB[IRBSL_SIDR_VBN] = .IRAB[IRBSL_FIRST_VBN];
791 0854 2 !RAB[IRBSW_SIDR_ID] = .IRAB[IRBSW_FIRST_ID];
792 0855 2
793 0856 2 ! Move the RFA of the current primary data record from its temporary
794 0857 2 location into the local NRP context and make it both the current primary
795 0858 2 data record and the primary data record for NRP positioning.
796 0859 2
797 0860 2 !RAB[IRBSL_UDR_VBN] = .IRAB[IRBSL_NEXT_VBN];
798 0861 2 !RAB[IRBSW_UDR_ID] = .IRAB[IRBSW_NEXT_ID];
799 0862 2
800 0863 2 !RAB[IRBSL_POS_VBN] = .IRAB[IRBSL_NEXT_VBN];
801 0864 2 !RAB[IRBSW_POS_ID] = .IRAB[IRBSW_NEXT_ID];
802 0865 2
803 0866 2 ! Setup up the key of reference of the current record, and move the key of
804 0867 2 the current record into keybuffer 1.
805 0868 2
806 0869 2 !RAB[IRBSB_CUR_KREF] = .IRAB[IRBSB_RP_KREF];
807 0870 2 CH$MOVE(.IFAB[IFBSW_KBUFSZ], KEYBUF_ADDR(2), KEYBUF_ADDR(1));
808 0871 2 RETURN;
809 0872 2
810 0873 1 END;
```

```
:TITLE RM3GET
:INT \V04-000\

.EXTRN RM$COMPARE_KEY, RMSFIND_BY_RRV
```

```

.EXTRN RMSKEY DESC, RMSLOCK
.EXTRN RMSNOREAD_LONG, RMSNOWRT_LONG
.EXTRN RMSPCKDEC_CHECK
.EXTRN RMSPOS KEY, RMSPOS RFA
.EXTRN RMSPOS SEQ, RMSQUERY LCK
.EXTRN RMSRECORD_ID, RMSRECORD KEY
.EXTRN RMSRECORD_VBN, RMSREC 0VHD
.EXTRN RMSRLSBKT, RMSRLU RECLAIM
.EXTRN RMSUNLOCK, RMSUNPACK_REC

.PSECT RMSRMS3,NOWRT, GBL, PIC,2

```

04	A9	02	8A	00000	SETUP_NRP_DATA:			
00A8	C9	70	A9	00 0004		B1CB2	#2, 4(IRAB)	: 0830
		00C2	C9	95 0000A		MOVL	112(IRAB), 168(IRAB)	: 0834
			08	12 0000E		TSTB	194(IRAB)	: 0839
00B8	C9	74	A9	80 00010		BNEQ	1\$	
				13 11 00016		MOVW	116(IRAB), 184(IRAB)	: 0841
00C0	C9	76	A9	80 00018	1\$:	BRB	2\$	
00B4	C9	7C	A9	00 001E		MOVW	118(IRAB), 192(IRAB)	: 0851
00BE	C9	00B2	C9	80 00024		MOVL	124(IRAB), 180(IRAB)	: 0852
00B0	C9	78	A9	00 002B	2\$:	MOVW	130(IRAB), 190(IRAB)	: 0853
00BC	C9	00B0	C9	80 00031		MOVL	120(IRAB), 176(IRAB)	: 0860
00AC	C9	78	A9	00 0038		MOVW	128(IRAB), 188(IRAB)	: 0861
00BA	C9	00B0	C9	80 0003E		MOVL	120(IRAB), 172(IRAB)	: 0863
00C3	C9	00C2	C9	90 00045		MOVW	128(IRAB), 186(IRAB)	: 0864
		50	00B4	CA 3C 0004C		MOVB	194(IRAB), 195(IRAB)	: 0869
		50	60	A9 C0 00051		MOVZWL	180(IFAB), R0	: 0870
60	B9	60	00B4	CA 28 00055		ADDL2	96(IRAB), R0	
				05 0005C		MOV3	180(IFAB), (R0), 296(IRAB)	
						RSB		: 0873

: Routine Size: 93 bytes, Routine Base: RMSRMS3 + 0000

: 811 0874 1

813
814 0875 1 %SBTTL 'RMSGET3B'
815 0876 1 GLOBAL ROUTINE RMSGET3B : RL\$RABREG =
816 0877 1
817 0878 1 ++
818 0879 1
819 0880 1 FUNCTIONAL DESCRIPTION:
820 0882 1 This routine implements the get/find operation for the
821 0883 1 indexed file organization.
822 0884 1
823 0885 1 CALLING SEQUENCE:
824 0886 1
825 0887 1 RMSGET3()
826 0888 1
827 0889 1 INPUT PARAMETERS:
828 0890 1
829 0891 1 R11 Impure area pointer
830 0892 1 R10 IFAB -- Pointer to IFAB
831 0893 1 R9 IRAB -- Pointer to IRAB
832 0894 1 R8 RAB -- pointer to users RAB
833 0895 1 ROP field options (NLK,ULK,RLK,LOC,NXR)
834 0896 1 RAC field = (SEQ, or KEY, or RFA)
835 0897 1 RFA field if RAC = RFA
836 0898 1 KBF,KSZ,KRF if RAC = KEY and KBF,KSZ if RAC = SEQ and LIM set
837 0899 1 UBF,USZ -- if a GET
838 0900 1
839 0901 1 IMPLICIT INPUTS:
840 0902 1
841 0903 1 IRAB fields:
842 0904 1
843 0905 1 IRBSV_UNLOCK_RP - current record should be unlocked before
844 0906 1 accessing new record.
845 0907 1 IRBSV_FIND_LAST - last operation was a FIND.
846 0908 1 IRBSV_SKIP_NEXT - last operation was sequential, the record
847 0909 1 described by the npn info is to be skipped
848 0910 1 and the record beyond it becomes the new
849 0911 1 record.
850 0912 1 IRBSL_KEYBUF (key buffer 1 or 2, maybe 3)
851 0913 1
852 0914 1 IFAB fields:
853 0915 1
854 0916 1 IFBSB_PLG_VER
855 0917 1 IFBSV_RU_RLK - if set, perform pseudo record locking
856 0918 1 IFBSV_RUP - if set, Recovery Unit is in progress
857 0919 1 IFBSW_KBUFSZ - size of each keybuffer
858 0920 1 IFBSV_NORECLK - record locking not required, i.e., not
859 0921 1 sharing the file and single stream only.
860 0922 1 IFBSV_WRTACC - if accessed for other than read only.
861 0923 1
862 0924 1 OUTPUT PARAMETERS:
863 0925 1
864 0926 1 RAB fields:
865 0927 1
866 0928 1 RFA of record found
867 0929 1 STV if io errors
868 0930 1 RBF,RSZ -- if a GET
869 0931 1

```
870 0932 1 IMPLICIT OUTPUTS:  
871 0933 1  
872 0934 1 IRAB fields:  
873 0935 1  
874 0936 1 IRBSV_UNLOCK RP  
875 0937 1 IRBSV_FIND LAST  
876 0938 1 IRBSV_SKIP NEXT  
877 0939 1 IRBSB_RP KREF  
878 0940 1 IRBSB_CUR KREF  
879 0941 1 IRBSL_KEYBUF (key buffer 1 or 2, maybe 3)  
880 0942 1 IRBSL_RBF User buffer address and size  
881 0943 1 IRBSW_RSZ  
882 0944 1  
883 0945 1 ROUTINE VALUE:  
884 0946 1  
885 0947 1 Internal RMS status code  
886 0948 1  
887 0949 1 SIDE EFFECTS:  
888 0950 1  
889 0951 1 Retrieved record maybe locked, and next record context is modified.  
890 0952 1  
891 0953 1 --  
892 0954 1  
893 0955 2 BEGIN  
894 0956 2  
895 0957 2 BUILTIN  
896 0958 2 AP,  
897 0959 2 TESTBITSC;  
898 0960 2  
899 0961 2 EXTERNAL REGISTER  
900 0962 2 COMMON_RAB_STR;  
901 0963 2  
902 0964 2 GLOBAL REGISTER  
903 0965 2 R_REC_ADDR_STR,  
904 0966 2 R_IDX_DFN_STR;  
905 0967 2  
906 0968 2 LOCAL  
907 0969 2 FLAGS : BLOCK[1],  
908 0970 2 STATUS;  
909 0971 2  
910 0972 2 MACRO  
911 0973 2 AP_STATUS = 0,0,1,0 %.  
912 0974 2 OK_WAT_STATUS = 0,1,1,0 %:  
913 0975 2 RU_DEL_STATUS = 0,2,1,0 %:  
914 0976 2  
915 0977 2  
916 0978 2 ! Continue to attempt to get the next record under the following  
917 0979 2 circumstances:  
918 0980 2  
919 0981 2 1. The status returned from GET_RECORD indicates the next record has  
920 0982 2 been deleted and RMS's access mode is sequential or random by key  
921 0983 2 value. Any key of reference. This can only happen if RMS has had  
922 0984 2 to wait for a record lock to be granted (status returned will be  
923 0985 2 RMS$_DEL), or RMS has positioned to and managed to lock a primary  
924 0986 2 data record that is marked RU_DELETE (status returned will be 0).  
925 0987 2  
926 0988 2 2. The status returned from GET_RECORD is an alternate success status
```

```
927 0989 2 | (OK_WAT), and the key of reference is NOT the primary key.
928 0990 2
929 0991 2
930 0992 2 | 3. An IDX_DFN has been allocated for the key descriptor, indicating
931 0993 2 | there are no problems with the key of reference.
932 0994 2
933 0995 2
934 0996 2 | Force the key descriptor to be initially zero.
935 0997 2
936 0998 2
937 0999 2
938 1000 2
939 1001 3
940 1002 3
941 1003 3
942 1004 3
943 1005 3
944 1006 3
945 1007 3
946 1008 3
947 1009 3
948 1010 4
949 1011 3
950 1012 3
951 1013 3
952 1014 3
953 1015 3
954 1016 3
955 1017 3
956 1018 3
957 1019 3
958 1020 3
959 1021 3
960 1022 6
961 1023 5
962 1024 5
963 1025 4
964 1026 4
965 1027 3
966 1028 4
967 1029 4
968 1030 4
969 1031 3
970 1032 3
971 1033 3
972 1034 3
973 1035 3
974 1036 3
975 1037 3
976 1038 3
977 1039 3
978 1040 3
979 1041 3
980 1042 3
981 1043 3
982 1044 3
983 1045 3

| (OK_WAT), and the key of reference is NOT the primary key.
| 3. An IDX_DFN has been allocated for the key descriptor, indicating
| there are no problems with the key of reference.
| Force the key descriptor to be initially zero.
| IDX_DFN = 0;
| FLAGS = 0;
WHILE 1
DO
  BEGIN
    STATUS = GET_RECORD (.FLAGS[OK_WAT_STATUS]
                         OR
                         TESTBITS(.FLAGS[RU_DEL_STATUS]));
    |
    | If an unqualified success, avoid the contorted mass of
    | logic below and exit immediately...
    IF .STATUS<0,16> EQU RMSSUC()
    THEN
      EXITLOOP;
    |
    | Check the key descriptor after return from GET_RECORD.
    | If still zero, something wrong with the key of
    | reference, so exit loop.
    IF .IDX_DFN EQU 0
    THEN
      EXITLOOP;
    |
    IF NOT ((.STATUS<0,16> EQU RMSERR(DEL)
              OR
              .STATUS EQU 0)
            AND
            .RAB[RAB$B_RAC] NEQU RAB$C_RFA)
    THEN
      IF NOT (.IDX_DFN[IDX$B_KEYREF] NEQU 0
              AND
              .STATUS<0,16> EQU RMSSUC(OK_WAT))
      THEN
        EXITLOOP;
    |
    | Let us back off from the radical position above concerning the
    | ambitious attempts to continue to get a record. We should
    | never go back for a record if all of the following are
    | true:
    |   - status from GET_RECORD is RMS$_DEL
    |   - primary key of reference
    |   - random access (keyed OR RFA access)
    |   - no dups on primary key
    |   - exact key match
    |
    | Under these circumstances, it is at least useless to
    | go back after a record, and sometimes downright WRONG!
```

```
984 1046 3
985 1047 3
986 1048 5
987 1049 4
988 1050 4
989 1051 4
990 1052 4
991 1053 4
992 1054 4
993 1055 4
994 1056 4
995 1057 3
996 1058 4
997 1059 4
998 1060 4
999 1061 3
1000 1062 3
1001 1063 3
1002 1064 3
1003 1065 3
1004 1066 3
1005 1067 3
1006 1068 3
1007 1069 3
1008 1070 3
1009 1071 3
1010 1072 3
1011 1073 3
1012 1074 3
1013 1075 3
1014 1076 3
1015 1077 2
1016 1078 2
1017 1079 2
1018 1080 2
1019 1081 2
1020 1082 2
1021 1083 2
1022 1084 2
1023 1085 2
1024 1086 2
1025 1087 2
1026 1088 2
1027 1089 2
1028 1090 2
1029 1091 2
1030 1092 3
1031 1093 2
1032 1094 2
1033 1095 2
1034 1096 2
1035 1097 2
1036 1098 2
1037 1099 2
1038 1100 2
1039 1101 2
1040 1102 2

: !  
: IF (.STATUS<0,16> EQLU RMSERR(DEL)  
:     AND  
:     .IDX_DFN[IDX$B_KEYREF] EQLU 0  
:     AND  
:     .RAB[RAB$B_RAC] NEQU RAB$C_SEQ  
:     AND  
:     NOT .IDX_DFN[IDX$V_DUPKEYS]  
:     AND  
:     NOT (.RAB[RAB$V_KGE] OR .RAB[RAB$V_KGT]))  
: THEN  
:     BEGIN  
:     STATUS = RMSERR(RNF);  
:     EXITLOOP;  
: END;  
:  
: If RMS has to perform a re-positioning, then either it had to have  
: waited for a record lock, or it positioned to a RU_DELETE marked  
: primary data record. Therefore, set the appropriate state bit either  
: of which will cause the correct lock to be released during  
: re-positioning, and so that the proper status will be returned if  
: RMS is able to position to a record.  
:  
: IF .STATUS NEQU 0  
: THEN  
:     FLAGS[OK_WAT_STATUS] = 1  
: ELSE  
:     FLAGS[RU_DEL_STATUS] = 1;  
:  
: CHSMOVE (.IDX_DFN[IDX$B_KEYSZ], KEYBUF_ADDR(1), KEYBUF_ADDR(2));  
: END;  
:  
: NOTE: AP is 0 if no special action is needed to unlock the RP on errors  
: produced due to user buffer/size errors. Otherwise it is 1. Save this  
: status.  
:  
: FLAGS[AP_STATUS] = .AP<0,1>;  
:  
: If RMS was successful at obtaining the next record, but at some earlier  
: time was forced to do a re-positioning because it had to wait for a record  
: lock, then change the status to an OK_WAT success. This can only happen  
: when RMS is positioning by means of an alternate index.  
:  
: IF .FLAGS[OK_WAT_STATUS]  
:     AND  
:     .STATUS<0,16> EQLU RMSSUC()  
: THEN  
:     STATUS = RMSSUC(OK_WAT);  
:  
: IRAB[IRBS$V_FIND_LAST] = 0;  
:  
: ! Obtain user buffer address and size for later probe.  
:  
: IRAB[IRBS$L_RBF] = .RAB[RAB$L_RBF];  
: IRAB[IRBS$W_RSZ] = .RAB[RAB$W_RSZ];
```

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1100
1101
1102
1103 2
1104 2
1105 2
1106 2
1107 2
1108 2
1109 3
1110 2
1111 2
1112 2
1113 3
1114 3
1115 3
1116 3
1117 3
1118 3
1119 3
1120 3
1121 3
1122 3
1123 3
1124 3
1125 3
1126 3
1127 4
1128 3
1129 3
1130 3
1131 4
1132 4
1133 4
1134 4
1135 4
1136 4
1137 4
1138 5
1139 3
1140 3
1141 3
1142 3
1143 3
1144 3
1145 3
1146 3
1147 4
1148 4
1149 4
1150 4
1151 4
1152 4
1153 4
1154 4
1155 4
1156 4
1157 3
1158 3
1159 2

! If the user has set the RAB\$V_LIM bit in the ROP field on a sequential
\$GET/\$FIND RMS reports whether the specified key exceeds the key of the
record found.

IF .STATUS
AND
.RAB[RAB\$B_RAC] EQL RAB\$C_SEQ)
AND
.RAB[RAB\$V_LIM]
THEN
BEGIN
LOCAL
KBF_ADDR : LONG,
KEYSIZE;
KEYSIZE = .RAB[RAB\$B_KSZ];
IF .KEYSIZE EQL 0
THEN
IF .IDX_DFN[IDX\$B_DATATYPE] EQL IDX\$C_STRING
OR .IDX_DFN[IDX\$B_SEGMENTS] GTR 1
THEN STATUS = RMSERR(KSZ)
ELSE KEYSIZE = .IDX_DFN[IDX\$B_KEYSZ];
BEGIN
MAP
KEYSIZE : BYTE;
IF .KEYSIZE GTRU .IDX_DFN[IDX\$B_KEYSZ]
THEN STATUS = RMSERR(KSZ)
END;
KBF_ADDR = .RAB[RAB\$L_KBF];
IFNORD(KEYSIZE, .KBF_ADDR, .RAB[RAB\$B_MODE],
STATUS = RMSERR(RBF));
IF .STATUS
THEN
BEGIN
AP = 3; ! Contiguous key compare
! The key of the current record has been previously saved in
keybuffer 2, and may now be used to determine whether the user
specified key limit has been exceeded.
IF RMSCOMPARE_KEY (KEYBUF_ADDR(2), .KBF_ADDR, .KEYSIZE) LSS 0
THEN STATUS = RMSSUC(OK_LIM);
END;
END;

```
; 1098
; 1099
; 1100
; 1101
; 1102
; 1103
; 1104
; 1105
; 1106
; 1107
; 1108
; 1109
; 1110
; 1111
; 1112
; 1113
; 1114
; 1115
; 1116
; 1117
; 1118
; 1119
; 1120
; 1121
; 1122
; 1123
; 1124
; 1125
; 1126
; 1127
; 1128
; 1129
; 1130
; 1131
; 1132
; 1133
; 1134
; 1135
; 1136
; 1137
; 1138
; 1139
; 1140
; 1141
; 1142
; 1143
; 1144
; 1145
; 1146
; 1147
; 1148
; 1149
; 1150
; 1151
; 1152
; 1153
; 1154
1160 2
1161 2
1162 2
1163 3
1164 3
1165 3
1166 3
1167 3
1168 3
1169 3
1170 3
1171 4
1172 4
1173 4
1174 4
1175 4
1176 5
1177 4
1178 4
1179 4
1180 4
1181 4
1182 4
1183 4
1184 5
1185 5
1186 5
1187 4
1188 4
1189 4
1190 3
1191 3
1192 3
1193 3
1194 4
1195 4
1196 4
1197 4
1198 4
1199 4
1200 4
1201 4
1202 4
1203 5
1204 5
1205 5
1206 5
1207 5
1208 5
1209 5
1210 5
1211 5
1212 5
1213 5
1214 5
1215 5
1216 5

IF .STATUS
THEN
BEGIN
IF .IRAB[IRBSV_FIND]
THEN
| This is a find operation don't
| return record etc.
BEGIN
IRAB[IRBSV_FIND_LAST] = 1;
| Set up the next record context for non-random $IND operations.
IF (.RAB[RAB$B_RAC] EQL RAB$C_SEQ)
THEN
SETUP_NRP_DATA()
| If this is a random $IND operation then save the RFA of the
| found primary data record as the current primary data record.
ELSE
BEGIN
IRAB[IRBSL_UDR_VBN] = .IRAB[IRBSL_NEXT_VBN];
IRAB[IRBSW_UDR_ID] = .IRAB[IRBSW_NEXT_ID];
END;
END
ELSE
| Return the user the data on the record.
BEGIN
LOCAL
RSZ : WORD;
RMSKEY_DESC(0);
| Add record overhead, and calculate record's size.
BEGIN
LOCAL
REC_SIZE,
RECORD_OVHD;
RECORD_OVHD = RMSREC_OVHD(0; REC_SIZE);
| If this primary data record had been updated within a Recovery
| Unit then retrieve its true size from the last two bytes of the
| reserved space.
IF .REC_ADDR[IRC$V_RU_UPDATE]
THEN
```

```
1155      1217 6      RSZ = .REC_ADDR + .RECORD_OVHD
1156      1218 6      + REC_SIZE
1157      1219 5      - IRC$C_DATSZFLD)<0,16>
1158      1220 5
1159      1221 5      ELSE RSZ = .REC_SIZE;
1160      1222 5
1161      1223 5      REC_ADDR = .REC_ADDR + .RECORD_OVHD;
1162      1224 4      END;
1163      1225 4
1164      1226 4      ! It will only be necessary to unpack the primary data record
1165      1227 4      if its a prologue 3 file.
1166      1228 4
1167      1229 5      IF (.IFAB[IFB$B_PLG_VER] EQLU PLG$C_VER_3)
1168      1230 4      THEN BEGIN
1169      1231 5      GLOBAL REGISTER
1170      1232 5      R_BKT_ADDR;
1171      1233 5
1172      1234 5
1173      1235 5
1174      1236 5
1175      1237 5      ! If the key of reference is the primary key then signal
1176      1238 5      RMSUNPACK_REC that the primary key for this data record maybe
1177      1239 5      found in expanded form within keybuffer 2.
1178      1240 5      IF .IRAB[IRBSB_RP_KREF] EQLU 0
1179      1241 5      THEN AP = 2
1180      1242 5      ELSE AP = 0;
1181      1243 5
1182      1244 5
1183      1245 5
1184      1246 5      RSZ = RMSUNPACK_REC(.IRAB[IRBSL_RECBUF],.RSZ);
1185      1247 4      END;
1186      1248 4
1187      1249 6      IF (.RAB[RAB$V_LOC] AND NOT (.IFAB[IFB$V_UPD])
1188      1250 5      AND
1189      1251 5      NOT (.BBLOCK[.IRAB[IRBSL_CURBDB], BDB$V_NOLOCATE]))
1190      1252 4      THEN BEGIN
1191      1253 5      ! We can do locate mode get
1192      1254 5      IF .IFAB[IFB$B_PLG_VER] LSSU PLG$C_VER_3
1193      1255 5      THEN RAB[RAB$L_RBF] = .REC_ADDR
1194      1256 5      ELSE RAB[RAB$L_RBF] = .IRAB[IRBSL_RECBUF];
1195      1257 5
1196      1258 5      RAB[RAB$W_RSZ] = .RSZ;
1197      1259 5
1198      1260 5      END
1199      1261 4      ELSE BEGIN
1200      1262 5      ! We must do move mode get
1201      1263 5
1202      1264 5
1203      1265 5      LOCAL
1204      1266 5      USZ : WORD;
1205      1267 5
1206      1268 5
1207      1269 5      IF .USZ EQL 0
1208      1270 5      THEN BEGIN
1209      1271 6
1210      1272 6
1211      1273 6      IF .FLAGS[AP_STATUS]
```

```
1212 1274 6
1213 1275 6
1214 1276 6
1215 1277 6
1216 1278 6
1217 1279 5
1218 1280 6
1219 1281 6
1220 1282 6
1221 1283 6
1222 1284 6
1223 1285 6
1224 1286 6
1225 1287 6
1226 1288 7
1227 1289 7
1228 1290 7
1229 1291 7
1230 1292 6
1231 1293 6
1232 1294 6
1233 1295 6
1234 1296 7
1235 1297 7
1236 1298 7
1237 1299 7
1238 1300 7
1239 1301 7
1240 1302 7
1241 1303 7
1242 1304 6
1243 1305 7
1244 1306 7
1245 1307 7
1246 1308 7
1247 1309 7
1248 1310 7
1249 1311 7
1250 1312 7
1251 1313 7
1252 1314 6
1253 1315 6
1254 1316 5
1255 1317 5
1256 1318 4
1257 1319 4
1258 1320 3
1259 1321 3
1260 1322 2
1261 1323 2
1262 1324 3
1263 1325 3
1264 1326 3
1265 1327 2
1266 1328 3
1267 1329 3
1268 1330 3

      THEN
          IRAB[IRBSV_UNLOCK_RP] = 1;      ! flag unlock RP
      STATUS = RMSERR(USZ);
      END
  ELSE
      BEGIN
          LOCAL
              UBF_ADDR;
          UBF_ADDR = .RAB[RABSL_UBF];
          IF .RSZ GTRU .USZ
          THEN
              BEGIN
                  RAB[RABSL_STV] = .RSZ;
                  RSZ = .USZ;
                  STATUS = RMSERR(RTB);
                  END;
          IF RMSNOWRT_LONG(.RSZ, .UBF_ADDR, .IRAB[IRBSB_MODE])
          THEN
              BEGIN
                  IF .FLAGS[AP_STATUS]
                  THEN
                      IRAB[IRBSV_UNLOCK_RP] = 1; ! flag unlock RP
                  STATUS = RMSERR(UBF);
                  END
              ELSE
                  BEGIN
                      RAB[RABSW_RSZ] = .RSZ;
                      RAB[RABSL_RBF] = .UBF_ADDR;
                      IF .IFABL[IFBSB_PLG_VER] LSSU PLGSC_VER_3
                      THEN
                          CH$MOVE(.RSZ, .REC_ADDR, .UBF_ADDR)
                      ELSE
                          CH$MOVE(.RSZ, .IRAB[IRBSL_RECBUF], .UBF_ADDR );
                  END;
              END;
          END;
      END;
  END;
END;
IF (.STATUS
    OR
    (.STATUS<0, 16> EQL RMSERR(RTB)))
THEN
BEGIN
    IF NOT .IRAB[IRBSV_FIND]
```

```
: 1269 1331 3
: 1270 1332 4
: 1</1 1333 4
: 1272 1334 4
: 1273 1335 4
: 1274 1336 4
: 1275 1337 4
: 1276 1338 4
: 1277 1339 4
: 1278 1340 4
: 1279 1341 4
: 1280 1342 4
: 1281 1343 4
: 1282 1344 4
: 1283 1345 5
: 1284 1346 5
: 1285 1347 5
: 1286 1348 5
: 1287 1349 5
: 1288 1350 5
: 1289 1351 5
: 1290 1352 5
: 1291 1353 5
: 1292 1354 5
: 1293 1355 5
: 1294 1356 5
: 1295 1357 5
: 1296 1358 6
: 1297 1359 6
: 1298 1360 6
: 1299 1361 5
: 1300 1362 6
: 1301 1363 6
: 1302 1364 6
: 1303 1365 6
: 1304 1366 6
: 1305 1367 6
: 1306 1368 5
: 1307 1369 5
: 1308 1370 4
: 1309 1371 3
: 1310 1372 3
: 1311 1373 3
: 1312 1374 2
: 1313 1375 3
: 1314 1376 3
: 1315 1377 3
: 1316 1378 3
: 1317 1379 3
: 1318 1380 3
: 1319 1381 3
: 1320 1382 3
: 1321 1383 3
: 1322 1384 3
: 1323 1385 3
: 1324 1386 3
: 1325 1387 3

      THEN
        BEGIN
          ! Set up the next record context for all $GET operations.
          SETUP_NRP_DATA();

          ! If dupes aren't allowed on primary key, save primary key value of
          ! this record in keybuffer 3 so reaccess and record lock window can
          ! be avoided on subsequent random FIND on primary key (cobol does
          ! this a lot).

          IF NOT .IDX_DFN[IDX$V_DUPKEYS]
          THEN
            BEGIN
              LOCAL
                TMP_REC_ADDR;
              TMP_REC_ADDR = .REC_ADDR;
              REC_ADDR = .RAB[RAB$L_RBF];
              IF RMSNOREAD LONG (.RAB[RAB$W_RSZ], .REC_ADDR,
                                 .IRAB[IRB$B_MODE])
              THEN
                STATUS = RMSERR(RBF);
                AP = 3; ! no overhead / expanded
                IF (.STATUS
                    OR
                    (.STATUS<0, 16> EQL RMSERR(RTB)))
              THEN
                BEGIN
                  GLOBAL REGISTER
                    R_BDB;
                  RMSRECORD_KEY (KEYBUF_ADDR(3));
                END;
                REC_ADDR = .TMP_REC_ADDR
              END;
            END;
          END;
        END
      ELSE
        BEGIN
          ! UNLOCK_RP is used as a flag on error conditions to indicate whether
          ! the record described by the current record (rp) information is to be
          ! unlocked or not. This will be the case when buffer errors are
          ! discovered after the new record has been locked, or if the current
          ! record before this operation was not unlocked at the beginning of
          ! this operation. In both cases the RFA of the record to unlock will be
          ! found in IRB$L_NEXT_VBN and IRB$W_NEXT_ID. In the former case,
          ! because these fields contain the RFA of the 'next' record, and in the
          ! latter, because the only reason why the current record was not
          ! unlocked at the beginning of the operation because it itself was
          ! being retrieved by the operation, and if this was the case then
        END
      END
    END
  END
END
```

```

1326 1388 3 | these same fields would contain the RFA of the current record from
1327 1389 3 | when it had been originally locked.
1328 1390 3
1329 1391 3 IF TESTBITS(.IRAB[IRBSV_UNLOCK_RP])
1330 1392 3 THEN
1331 1393 3 RMSUNLOCK (.IRAB[IRBSL_NEXT_VBN], .IRAB[IRBSW_NEXT_ID]);
1332 1394 3
1333 1395 3 | If end-of-file has been reached, set the corresponding IRAB bit.
1334 1396 3
1335 1397 4 IF .STATUS EQL RMSERR(EOF)
1336 1398 3 THEN
1337 1399 3 IRAB[IRBSV_EOF] = 1;
1338 1400 3
1339 1401 3 | There is no longer a current primary data record.
1340 1402 3
1341 1403 3 IRAB[IRBSL_UDR_VBN] = 0;
1342 1404 3 IRAB[IRBSW_UDR_ID] = 0;
1343 1405 2 END;
1344 1406 2
1345 1407 2 RAB[RABSL_RFA0] = .IRAB[IRBSL_UDR_VBN];
1346 1408 2 RAB[RABSW_RFA4] = .IRAB[IRBSW_UDR_ID];
1347 1409 2
1348 1410 2 IRAB[IRBSV_FIND] = 0;
1349 1411 3 BEGIN
1350 1412 3
1351 1413 3 GLOBAL REGISTER
1352 1414 3 R_BDB;
1353 1415 3
1354 1416 3 IF (BDB = .IRAB[IRBSL_CURBDB]) NEQ 0
1355 1417 3 THEN
1356 1418 3 RMSRLSBKT(0);
1357 1419 3
1358 1420 3 IRAB[IRBSL_CURBDB] = 0;
1359 1421 2 END;
1360 1422 2 RETURN .STATUS
1361 1423 2
1362 1424 1 END;

```

00FC 8F BB 00000 RMSGET3B::									
						PUSHR	#^M<R2,R3,R4,R5,R6,R7>		0876
						SUBL2	#16, SP		1
						CLRL	IDX_DFN		1
						CLRL	FLAGS		1
						CLRL	R0		1
						BBCC	#2, FLAGS, 2\$		1
						INCL	R0		1
						EXTZV	#1, #1, FLAGS, R1		1
						BISL3	R1, R0, -(SP)		1
						BSBW	GET_RECORD		1
						ADDL2	#4, SP		1
						MOVL	R0, STATUS		1
						CMPW	STATUS, #1		1
						BEQL	10\$		1
51	04	AE	01	0000V	30 0001F				1010
		7E	50						
			5E	04	00 00022				
			6E	50	00 00025				
			01	6E	B1 00028				
				6C	13 0002B				

63	6E	85A4	05	18	000E9	BLEQU	15\$	1138	
	53	30	A8	D0	000FO	MOVZWL	#34212, STATUS	1141	
	50	0A	A9	0C	000F4	MOVL	48(RAB\$), KBF_ADDR	1143	
	6E	858C	05	12	000F9	PROBER	10(IRAB\$), KEYSIZE, (KBF_ADDR)		
	18		6E	E9	00100	BNEQ	16\$		
	5C		03	D0	00103	MOVZWL	#34188, STATUS	1145	
	51	00B4	CA	3C	00106	BLBC	STATUS, 17\$	1148	
	51	60	A9	C0	0010B	MOVL	#3, AP	1154	
			0000G	30	0010F	MOVZWL	180(IFAB), R1		
				50	00112	ADDL2	96(IRAB), R1		
				05	18	BSBW	RMS_COMPARE_KEY		
				50	00114	TSTL	R0		
	6E	8051	8F	3C	00116	BGEQ	17\$		
	03		6E	E8	0011B	MOVZWL	#32849, STATUS	1156	
		010A	31	0011E	BLBS	STATUS, 18\$	1161		
1D	05	A9	01	E1	00121	BRW	38\$		
	04	A9	20	88	00126	BBC	#1, 5(IRAB), 21\$	1165	
		1E	A8	95	0012A	BISB2	#32, 4(IRAB)	1172	
			05	12	0012D	TSTB	30(RAB)	1176	
			FE71	30	0012F	BNEQ	19\$		
			0D	11	00132	BSBW	SETUP_NRP_DATA	1178	
	00B0	C9	78	A9	D0	BRB	20\$		
	008C	C9	0080	C9	80	MOVW	120(IRAB), 176(IRAB)	1185	
				72	11	00134	0D	128(IRAB), 188(IRAB)	1186
				7E	D4	00141	BRB	29\$	1165
				0000G	30	00143	CLRL	-(SP)	1199
		5E		04	C0	00145	BSBW	RMSKEY_DESC	
				51	D4	00148	ADDL2	#4, SP	
				0000G	30	0014D	CLRL	R1	1209
0E	66		06	E1	00150	BSBW	RMSREC_OVHD		
	56		50	C1	00154	BBC	#6, (REC_ADDR), 22\$	1215	
		FE	A142	9F	00158	ADDL3	RECORD_OVHD, REC_ADDR, R2	1217	
	08	AE	9E	B0	0015C	PUSHAB	-2(REC_SIZE)[R2]		
			04	11	00160	MOVW	0(SP)+, RSZ		
	08	AE	51	B0	00162	BRB	23\$	1221	
	56		50	C0	00166	22\$:	REC_SIZE, RSZ	1223	
	03	00B7	CA	91	00169	ADDL2	RECORD_OVHD, REC_ADDR	1229	
			1C	12	0016E	CMPB	183(IFAB), #3		
		00C2	C9	95	00170	BNEQ	26\$		
			05	12	00174	TSTB	194(IRAB)	1240	
	5C		02	D0	00176	BNEQ	24\$		
			02	11	00179	MOVL	#2, AP	1242	
			5C	D4	0017B	BRB	25\$		
	51	08	AE	3C	0017D	24\$:	CLRL	AP	1244
	50	68	A9	D0	00181	MOVZWL	RSZ, R1	1246	
			0000G	30	00185	MOVL	104(IRAB), R0		
	08	AE	50	B0	00188	BSBW	RMSUNPACK_REC		
	27	06	A8	E9	0018C	MOVW	R0, RSZ		
22	22	AA	03	E0	00190	BLBC	6(RAB), 30\$	1249	
	50	20	A9	D0	00195	BBS	#3, 34(IFAB), 30\$		
	0A	A0	04	E0	00199	MOVL	32(IRAB), R0	1251	
	03	00B7	CA	91	0019E	BBS	#4, 10(R0), 30\$	1254	
	28	A8	06	1E	001A3	CMPB	183(IFAB), #3		
			56	D0	001A5	BGEQ	27\$		
	28	A8	05	11	001A9	MOVL	REC_ADDR, 40(RAB)	1256	
		68	A9	D0	001AB	BRB	28\$		
						MOVL	104(IRAB), 40(RAB)	1258	

22	A8	08	AE	80 00180	28\$:	MOVW	RSZ, 34(RAB)	1259	
	50	20	A8	80 00185	29\$:	BRB	37\$	1249	
	04	04	OF	12 00188	30\$:	MOVW	32(RAB), USZ	1267	
05	A9	04	AE	E9 001BD		BNEQ	32\$	1269	
	6E	86F4	20	88 001C1	31\$:	BLBC	FLAGS, 31\$	1273	
			5C	11 001CA		BISB2	#32 5(IRAB)	1275	
0C	AE	24	A8	DO 001CC	32\$:	MOVZWL	#34548, STATUS	1277	
	50	08	AE	B1 001D1		MOVL	36(RAB), UBF_ADDR	1269	
			OE	1B 001D5		CMPW	RSZ, USZ	1284	
0C	A8	08	AE	3C 001D7		BLEQU	33\$	1286	
08	AE	50	BO	001DC		MOVZWL	RSZ, 12(RAB)	1289	
	6E	81A8	8F	3C 001E0		MOVZWL	USZ, RSZ	1290	
	7E	0A	A9	9A 001E5	33\$:	MOVZBL	#33192, STATUS	1291	
		10	AE	DD 001E9		PUSHL	10(IRAB), -(SP)	1294	
	7E	10	AE	3C 001EC		MOVZWL	UBF_ADDR	1222	
			0000G	30 001F0		BSBW	RSZ, -(SP)	1222	
			OC	CO 001F3		ADDL2	RMSNOWRT_LONG	1222	
			50	E9 001F6		BLBC	#12, SP	1222	
		04	AE	E9 001F9		BLBC	RO, 35\$	1298	
05	A9	20	88 001FD		BISB2	FLAGS, 34\$	1300		
	6E	86EC	8F	3C 00201	34\$:	MOVZWL	#32 5(IRAB)	1302	
		20	11 00206		BRB	37\$	1294		
22	A8	08	AE	80 00208	35\$:	MOVW	RSZ, 34(RAB)	1306	
28	A8	0C	AE	DO 0020D		MOVL	UBF_ADDR, 40(RAB)	1307	
		03	00B7	CA 91 00212		CMPB	183(IFAB), #3	1308	
			08	1E 00217		BGEQU	36\$	1222	
0C	BE	66	08	AE 28 00219		MOV3	RSZ, (REC_ADDR), @UBF_ADDR	1310	
0C	BE	68	B9	08	AE 28 00221	36\$:	MOV3	RSZ, @104(IRAB), @UBF_ADDR	1312
		07		6E E8 00228	37\$:	BLBS	STATUS, 39\$	1324	
		81A8	8F	6E B1 00228	38\$:	CMPW	STATUS, #33192	1326	
			4C	12 00230		BNEQ	43\$	1222	
6D	05	A9	01	E0 00232	39\$:	BBS	#1, 5(IRAB), 46\$	1330	
			FD69	30 00237		BSBW	SETUP_NRP_DATA	1336	
		66	1C	A7 E8 0023A		BLBS	28(IDX_DFN), 46\$	1343	
		52		56 D0 0023E		MOVL	REC_ADDR, TMP_REC_ADDR	1350	
		56	28	A8 DO 00241		MOVL	40(RAB), REC_ADDR	1351	
		7E	0A	A9 9A 00245		MOVZBL	10(IRAB), -(SP)	1354	
			56	DD 00249		PUSHL	REC_ADDR	1353	
		7E	22	A8 3C 0024B		MOVZWL	34(RAB), -(SP)	1222	
			0000G	30 0024F		BSBW	RMSNOREAD_LONG	1222	
			OC	CO 00252		ADDL2	#12, SP	1222	
		5E	50	E9 00255		BLBC	RO, 40\$	1222	
		05	8654	8F 3C 00258		MOVZWL	#34388, STATUS	1356	
		5C	03	DO 0025D	40\$:	MOVL	#3, AP	1357	
		07	6E	E8 00260		BLBS	STATUS, 41\$	1358	
		81A8	8F	6E B1 00263		CMPW	STATUS, #33192	1360	
			0F	12 00268		BNEQ	42\$	1222	
			50	CA 3C 0026A	41\$:	MOVZWL	180(IFAB), RO	1367	
			60	B940 3F 0026F		PUSHW	@96(IRAB)[RO]	1222	
			0000G	30 00273		BSBW	RMSRECORD_KEY	1222	
			04	CO 00276		ADDL2	#4, SP	1222	
		5E	52	DO 00279	42\$:	MOVL	TMP_REC_ADDR, REC_ADDR	1369	
		56		26 11 0027C		BRB	46\$	1324	
0C	04	A9	0D	E5 0027E	43\$:	BBCC	#13, 4(IRAB), 44\$	1391	

52	0080	C9 3C 00283	MOVZWL 128(IRAB), R2	1393
51	78	A9 D0 00288	MOVL 120(IRAB), R1	1393
		0000G 30 0028C	BSBW RMSUNLOCK	1393
0000827A	8F	6E D1 0028F	CMPL STATUS, #33402	1397
		44\$:		1397
04	A9	04 12 00296	BNEQ 45\$	1399
		02 88 00298	BISB2 #2, 4(IRAB)	1399
10	A8	00B0 C9 D4 0029C	CLRL 176(IRAB)	1403
14	A8	00BC C9 B4 002A0	CLRW 188(IRAB)	1404
05	A9	00B0 C9 D0 002A4	MOVL 176(IRAB), 16(RAB)	1407
		46\$:	MOVW 188(IRAB), 20(RAB)	1408
		02 8A 002B0	BICB2 #2, 5(IRAB)	1410
	54	20 A9 D0 002B4	MOVL 32(IRAB), BDB	1416
		08 13 002B8	BEQL 47\$	1418
		7E D4 002BA	CLRL -(SP)	1418
		0000G 30 002BC	BSBW RMSRLSBKT	1420
5E		04 C0 002BF	ADDL2 #4, SP	1422
50		20 A9 D4 002C2	CLRL 32(IRAB)	1422
5E		8E D0 002C5	MOVL STATUS, R0	1424
		0C C0 002C8	ADDL2 #12, SP	1424
		00FC 8F BA 002CB	POPR #^M<R2,R3,R4,R5,R6,R7>	1424
		05 002CF	RSB	1424

: Routine Size: 720 bytes. Routine Base: RMS\$RMS3 + 005D

: 1363 1425 1

GET_RECORD

1365 1426 1 XSBTTL 'GET_RECORD'
1366 1427 1 ROUTINE GET_RECORD (REPOS_STATUS) : L_GET_RECORD =
1367 1428 1
1368 1429 1 :++
1369 1430 1
1370 1431 1 FUNCTIONAL DESCRIPTION:
1371 1432 1
1372 1433 1 This routine implements the actual retrieval of the
1373 1434 1 data record for internal RMS usage. The use request
1374 1435 1 is checked for valid input parameters and all internal state
1375 1436 1 information is setup to retrieve the record.
1376 1437 1 Then current record is unlocked if required and the
1377 1438 1 requested record retrieved and locked if required. All NRP
1378 1439 1 update data is saved in the IRAB but is not placed in the
1379 1440 1 NRP fields of the IRAB.
1380 1441 1
1381 1442 1 CALLING SEQUENCE:
1382 1443 1
1383 1444 1 GET_RECORD()
1384 1445 1
1385 1446 1 INPUT PARAMETERS:
1386 1447 1
1387 1448 1 REPOS_STATUS - if 1, then RMS is performing a re-positioning.
1388 1449 1
1389 1450 1 IMPLICIT INPUTS:
1390 1451 1
1391 1452 1 Same as for RMSGET3 or RM\$FIND3
1392 1453 1
1393 1454 1 OUTPUT PARAMETERS:
1394 1455 1
1395 1456 1 IRAB context setup for retrieved record:
1396 1457 1
1397 1458 1 CURBDB,RFA_VBN,RFA_ID,SAVE_POS,FIRST_VBN
1398 1459 1 FIRST_ID,REC_ADDR,NEXT_VBN,NEXT_ID
1399 1460 1
1400 1461 1 RAB The DCT field is cleared in all cases.
1401 1462 1
1402 1463 1 If the value of the routine is a success status then
1403 1464 1 the AP = 0 if no special action is needed to unlock the RP
1404 1465 1 and is 1 if special action is needed, on errors detected after
1405 1466 1 this routine.
1406 1467 1
1407 1468 1 IMPLICIT OUTPUTS:
1408 1469 1
1409 1470 1 IRBSV_UNLOCK RP
1410 1471 1 IRBSV_FIND LAST = 0
1411 1472 1 IRBSB_RP_KREF
1412 1473 1
1413 1474 1 ROUTINE VALUE:
1414 1475 1
1415 1476 1 Internal RMS status code
1416 1477 1
1417 1478 1 SIDE EFFECTS:
1418 1479 1
1419 1480 1 Retrieved record maybe locked.
1420 1481 1 Old current record may have been unlocked.
1421 1482 1 The data bucket for the retrieved record is accessed.

GET_RECORD

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1483 1 NLK is ignored if the process is in a Recovery Unit.
1484 1 If a primary data record is found to have been deleted within a Recovery
1485 1 Unit it might be deleted for good.
1486 1 If a primary data record is found to have been updated within a Recovery
1487 1 Unit it might be re-formatted.
1488 1 --
1489 1 BEGIN
1490 1 BUILTIN
1491 2 TESTBITCC.
1492 2 TESTBITSC.
1493 2 AP;
1494 2 EXTERNAL REGISTER
1495 2 R_REC_ADDR STR,
1496 2 R_IDX_DFN STR,
1497 2 COMMON_RAB_STR;
1498 2 LABEL
1499 2 UNLOCK,
1500 2 KEY;
1501 2
1502 2
1503 2
1504 2
1505 2
1506 2
1507 2 IRAB[IRBSB_CACHEFLGS] = 0;
1508 2 IRAB[IRBSW_SRCHFLAGS] = 0;
1509 2 IRAB[IRBSV_DUP] = 0;
1510 2
1511 2 ! Based on the record access mode (RAC) of this operation (GET/FIND)
1512 2 set up the IRAB RP fields, key buffer 2 etc to retrieve the record.
1513 2
1514 2
1515 2 RAB[RAB\$L_DCT] = 0;
1516 2
1517 2 ! Get record block 1 --- set up the IRAB search context data to get the
1518 2 record the user is requesting, and unlock the current record if this is
1519 2 required.
1520 2
1521 2 UNLOCK :
1522 3 BEGIN
1523 3
1524 3 CASE .RAB[RAB\$B_RAC] FROM RAB\$C_SEQ TO RAB\$C_RFA OF
1525 3 SET
1526 3
1527 3 [RAB\$C_SEQ] :
1528 3
1529 3 !+
1530 3 Sequential Access:
1531 3
1532 3 Setup to retrieve the record associated with the NRP if this
1533 3 is a GET and the last operation was not a FIND or if this is
1534 3 a FIND.
1535 3
1536 3 If last operation was a FIND and this operation is a GET then
1537 3 retrieve record which is described by the NRP if that FIND was
1538 3 sequential.
1539 3

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: 1479      1540 3
: 1480      1541 3
: 1481      1542 3
: 1482      1543 3
: 1483      1544 3
: 1484      1545 3
: 1485      1546 3
: 1486      1547 4
: 1487      1548 4
: 1488      1549 4
: 1489      THEN .IRAB[ IRBSV_CON_EOF ]
: 1490      1550 4
: 1491      1551 4
: 1492      1552 4
: 1493      1553 5
: 1494      1554 5
: 1495      1555 5
: 1496      THEN .IRAB[IRBSV_FIND]
: 1497      1556 4
: 1498      1557 5
: 1499      1558 5
: 1500      1559 5
: 1501      1560 5
: 1502      1561 5
: 1503      1562 5
: 1504      1563 5
: 1505      1564 5
: 1506      1565 5
: 1507      1566 5
: 1508      1567 5
: 1509      1568 5
: 1510      1569 5
: 1511      1570 5
: 1512      1571 5
: 1513      1572 6
: 1514      1573 5
: 1515      1574 5
: 1516      1575 5
: 1517      1576 5
: 1518      1577 5
: 1519      1578 5
: 1520      1579 5
: 1521      1580 5
: 1522      1581 5
: 1523      1582 5
: 1524      1583 6
: 1525      1584 6
: 1526      1585 6
: 1527      1586 5
: 1528      1587 5
: 1529      1588 5
: 1530      1589 5
: 1531      1590 5
: 1532      1591 5
: 1533      1592 5
: 1534      1593 5
: 1535      1594 5
: 1536      1595 5
: 1537      1596 5

      | If the FIND was random then change the NRP data for the record
      | which was found and retrieve it.

      | Note that a sequential FIND following a random FIND returns to
      | the sequential next record (NRP). That is to say that the
      | random operation will not change the NRP VBN and ID fields.

      - BEGIN

      IF .IRAB[ IRBSV_CON_EOF ]
      THEN RETURN RMSERR(EOF);

      IF (.IRAB[IRBSV_FIND_LAST]
      AND
      NOT (.IRAB[IRBSV_FIND]))
      THEN BEGIN

      | NOTE: keybuffer 2 contains key value, RP_KREF has key of
      | reference, RP_VBN and RP_ID contains record's RFA/RRV, and
      | SAVE_DUP contains the duplicate position count.

      IF TESTBITSC(IRAB[IRBSV_SKIP_NEXT])
      THEN

      | Last find was sequential so retrieve the record described
      | by the NRP unless RMS is already positioned at the end
      | of the file.

      IF .IRAB[IRBSV_EOF]
      THEN RETUF . RMSERR (EOF)
      ELSE RMSKEY_DESC(.IRAB[IRBSB_CUR_KREF])
      ELSE

      | Last operation was a find random and this operation is a
      | get sequential. Setup the local NRP context so that this
      | same record will be retrieved. Since there has been no
      | intervening operation, calling SETUP_NRP_DATA will
      | accomplish this.

      BEGIN
      SETUP_NRP_DATA();
      RMSKEY_DESC(.IRAB[IRBSB_RP_KREF]);
      END;

      IRAB[IRBSB_KEYSZ] = .IDX_DFN[IDXSB_KEYSZ];

      | Unless no lock is desired on this record (and the process
      | is not within a Recovery Unit), leave this block to avoid
      | unlocking the current record. This avoids a potential
      | window where the record is unlocked as it is reaccessed on
      | this get operation.

      IF NOT .RAB[RABSV_NLK]
```

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1536 1597 5
1537 1598 5
1538 1599 5
1539 1600 5
1540 1601 5
1541 1602 5
1542 1603 4
1543 1604 5
1544 1605 5
1545 1606 5
1546 1607 5
1547 1608 5
1548 1609 5
1549 1610 5
1550 1611 5
1551 1612 5
1552 1613 5
1553 1614 5
1554 1615 5
1555 1616 5
1556 1617 5
1557 1618 5
1558 1619 5
1559 1620 5
1560 1621 5
1561 1622 5
1562 1623 5
1563 1624 5
1564 1625 5
1565 1626 4
1566 1627 4
1567 1628 3
1568 1629 3
1569 1630 3
1570 1631 3 KEY :
1571 1632 4
1572 1633 4
1573 1634 4
1574 1635 4
1575 1636 4
1576 1637 4
1577 1638 4
1578 1639 4
1579 1640 4
1580 1641 4
1581 1642 4
1582 1643 4
1583 1644 4
1584 1645 4
1585 1646 4
1586 1647 4
1587 1648 4
1588 1649 4
1589 1650 5
1590 1651 5
1591 1652 7
1592 1653 6

      OR
      .IFAB[IFBSV_RUP]
      THEN
      LEAVE UNLOCK;
      END
      ELSE
      BEGIN
      ! Return immediately if already at end-of-file.
      IF .IRAB[IRBSV_EOF]
      THEN
      RETURN RMSERR(EOF);
      IRAB[IRBSV_SKIP_NEXT] = 1;
      ! First time call after $CONNECT or $REWIND. Then what we
      ! want to retrieve is the very first record, so don't skip
      ! next record.
      IF .IRAB[IRBSL_CUR_VBN] EQL 0
      THEN
      IRAB[IRBSV_SKIP_NEXT] = 0;
      RETURN_ON_ERROR (RMSKEY_DESC(.IRAB[IRBSB_CUR_KREF]));
      IRAB[IRBSB_KEYSZ] = .IDX_DFN[IDXSB_KEYSZ];
      CHSMOVE(.IDX_DFN[IDXSB_KEYSZ], KEYBUF_ADDR(1), KEYBUF_ADDR(2));
      END;
      END;
      [RABSC_KEY] :
      BEGIN
      LOCAL
      KEYSIZE : BYTE,
      KBF_ADDR : LONG;
      IRAB[IRBSV_CON_EOF] = 0;
      IRAB[IRBSV_SKIP_NEXT] = 0;
      RETURN_ON_ERROR (RMSKEY_DESC(.RAB[RABSB_KRF]));
      KEYSIZE = .RAB[RABSB_KSZ];
      ! Check and setup for user key value.
      IF .IDX_DFN[IDXSB_DATATYPE] EQL IDXSC_STRING
      OR .IDX_DFN[IDXSB_SEGMENTS] GTR 1
      THEN
      BEGIN
      IF ((.KEYSIZE EQL 0)
      OR
```

GET_RECORD

```

: 1593 1654 6
: 1594 1655 5
: 1595 1656 5
: 1596 1657 5
: 1597 1658 5
: 1598 1659 4
: 1599 1660 5
: 1600 1661 5
: 1601 1662 5
: 1602 1663 5
: 1603 1664 5
: 1604 1665 5
: 1605 1666 5
: 1606 1667 5
: 1607 1668 5
: 1608 1669 5
: 1609 1670 4
: 1610 1671 4
: 1611 1672 4
: 1612 1673 4
: 1613 P 1674 4
: 1614 1675 4
: 1615 1676 4
: 1616 1677 4
: 1617 1678 4
: 1618 1679 4
: 1619 1680 4
: 1620 1681 4
: 1621 1682 4
: 1622 1683 4
: 1623 1684 4
: 1624 1685 4
: 1625 1686 4
: 1626 1687 4
: 1627 1688 4
: 1628 1689 4
: 1629 1690 4
: 1630 1691 4
: 1631 1692 4
: 1632 1693 4
: 1633 1694 5
: 1634 1695 4
: 1635 1696 5
: 1636 1697 5
: 1637 1698 5
: 1638 1699 5
: 1639 1700 5
: 1640 1701 4
: 1641 1702 4
: 1642 1703 4
: 1643 1704 4
: 1644 1705 5
: 1645 1706 5
: 1646 1707 5
: 1647 1708 5
: 1648 1709 4
: 1649 1710 4

        (.KEYSIZE GTRU .IDX_DFN[IDX$B_KEYSZ]))
        THEN RETURN RMSERR(KSZ);

        ELSE END
        BEGIN
            IF .KEYSIZE EQL 0
            THEN KEYSIZE = .IDX_DFN[IDX$B_KEYSZ];
            IF .KEYSIZE NEQU .IDX_DFN[IDX$B_KEYSZ]
            THEN RETURN RMSERR(KSZ);
        END;

        IRAB[IRB$B_KEYSZ] = .KEYSIZE;
        KBF ADDR = .RAB[RAB$L_KBF];
        IFNORD(KEYSIZE, .KBF ADDR, IRAB[IRB$B_MODE]);
        RETURN RMSERR(KBF);

        ! Move the user's key into keybu.fier 2.
        CHSMOVE(.KEYSIZE, .KBF_ADDR, KEYBUF_ADDR(2));

        ! If key type is packed decimal then check it for valid nibbles.
        IF .IDX_DFN[IDX$B_DATATYPE] EQLU IDX$C_PACKED
        THEN RETURN_ON_ERROR (RMSPCKDEC_CHECK());

        ! Check that key match is logically consistent.
        IF .RAB[RAB$V_KGE]
        THEN
            IF .RAB[RAB$V_KGT]
            THEN RETURN RMSERR(POP)
            ELSE BEGIN
                IRAB[IRB$V_SRCHGE] = 1;
                LEAVE KEY
            END
        ELSE
            IF .RAB[RAB$V_KGT]
            THEN BEGIN
                IRAB[IRB$V_SRCHGT] = 1;
                LEAVE KEY
            END;

```

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1650 1711 4
1651 1712 4
1652 1713 4
1653 1714 4
1654 1715 4
1655 1716 4
1656 1717 4
1657 1718 4
1658 1719 4
1659 1720 4
1660 1721 4
1661 1722 4
1662 1723 4
1663 1724 4
1664 1725 4
1665 1726 4
1666 1727 4
1667 1728 4
1668 1729 4
1669 1730 4
1670 1731 4
1671 1732 4
1672 1733 4
1673 1734 4
1674 1735 4
1675 1736 4
1676 1737 4
1677 1738 4
1678 1739 4
1679 1740 4
1680 1741 5
1681 1742 5
1682 1743 5
1683 1744 5
1684 1745 5
1685 1746 5
1686 1747 5
1687 1748 5
1688 1749 5
1689 1750 5
1690 1751 5
1691 1752 4
1692 1753 5
1693 1754 5
1694 1755 5
1695 1756 5
1696 1757 5
1697 1758 5
1698 1759 5
1699 1760 5
1700 1761 5
1701 1762 5
1702 1763 5
1703 1764 5
1704 1765 5
1705 1766 5
1706 1767 5

+ At this point we have determined that this a random access for an
exact match by key. Now try to find out if this is for the
current record, i.e., is it the same one we just got. This will
be checked only for primary key. The following conditions must be
met to take this optimization:
Previous operation was a GET.
This operation is for primary key.
Duplicates aren't allowed on primary key.
The full key size is being used.
The key value matches the saved primary key value of the
current record (in keybuffer 3).
There is a current record (rp_vbn neq 0).
The current record is already locked, if locking required.
The new record is to be locked, if locking required.
- IF .IRAB[IRBSV_FIND_LAST]
  OR
    .IDX_DFN[IDX$B_KEYREF] NEQ 0
  OR
    .IDX_DFN[IDX$V_DUPKEYS]
THEN
  LEAVE KEY;

IF .IRAB[IRBSB_KEYSZ] NEQ .IDX_DFN[IDX$B_KEYSZ]
THEN
  LEAVE KEY;

BEGIN
  LOCAL
    SIZE;
  SIZE = .IRAB[IRBSB_KEYSZ];
  IF NOT CHSEQL(.SIZE, KEYBUF_ADDR(2), .SIZE, KEYBUF_ADDR(3))
  THEN
    LEAVE KEY;

END;
BEGIN
  LOCAL
    VBN;
  IF (VBN = .IRAB[IRBSL_UDR_VBN]) EQL 0
  THEN
    LEAVE KEY;

  ! If record locking is required, make sure this record is already
  ! locked, otherwise it may be deleted or locked by another
  ! accessor. Also that the new record is to be locked also,
  ! otherwise there is potentially an obscure window where it could
  ! be deleted while reaccessing the bucket after the current lock is
  ! released.

```



```
; 1764 1825 4      ! - exact match by primary key
; 1765 1826 4      - no dups
; 1766 1827 4
; 1767 1828 4      The above conditions will give us the first non-deleted record,
; 1768 1829 4      regardless of key (since we temporarily used sequential access).
; 1769 1830 4      And since KEYBUFFER 3 still has the key from when we last had
; 1770 1831 4      this record (before we released it), it will match our search
; 1771 1832 4      key. I have attempted in RMSGET3B to prevent coming back here a
; 1772 1833 4      second time if these conditions are true.
; 1773 1834 4
; 1774 1835 4      IRAB[IRBSV_DUP] = 1;
; 1775 1836 4      LEAVE UNLOCK;
; 1776 1837 4
; 1777 1838 3      END:                      ! of block KEY
; 1778 1839 3
; 1779 1840 3      [RAB$C_RFA] :
; 1780 1841 3
; 1781 1842 3      ! RFA access -- check RFA for legality and setup for primary key
; 1782 1843 3      access for next record pointer (NRP) data.
; 1783 1844 3
; 1784 1845 4      BEGIN
; 1785 1846 4
; 1786 1847 4      IRAB[IRBSV_CON_EOF] = 0;
; 1787 1848 4
; 1788 1849 4      IRAB[IRBSV_SKIP_NEXT] = 0;      ! flag random access
; 1789 1850 4
; 1790 1851 4      IF .RAB[RAB$L_RFA0] EQL 0
; 1791 1852 4      OR
; 1792 1853 4      .RAB[RAB$W_RFA4] EQL 0
; 1793 1854 4      THEN
; 1794 1855 4      RETURN RMSERR(RFA);
; 1795 1856 4
; 1796 1857 4      RETURN_ON_ERROR (RMSKEY_DESC(0));
; 1797 1858 4
; 1798 1859 3      END;
; 1799 1860 3
; 1800 1861 3      [OUTRANGE] :
; 1801 1862 3      RETURN RMSERR(RAC);
; 1802 1863 3
; 1803 1864 3      TES:
; 1804 1865 3
; 1805 1866 3      ! The current record is now unlocked before accessing the new record,
; 1806 1867 3      unless it has already been determined that the new record is the same as
; 1807 1868 3      the old current record, in which case this block was left and this code
; 1808 1869 3      is skipped.
; 1809 1870 3
; 1810 1871 3      IF TESTBITS(1RAB[IRBSV_UNLOCK_RP])
; 1811 1872 3      THEN
; 1812 1873 3
; 1813 1874 3      ! If RMS is performing a re-positioning then unlock the record
; 1814 1875 3      positioned to during the previous positioning attempt; otherwise,
; 1815 1876 3      unlock the current record (locked during the previous positioning
; 1816 1877 3      operation) if there is one.
; 1817 1878 3
; 1818 1879 3      IF .REPOS_STATUS
; 1819 1880 3      THEN
; 1820 1881 3      RMSUNLOCK (.IRAB[IRBSL_NEXT_VBN], .IRAB[IRBSW_NEXT_ID])
```


GET_RECORD

```
; 1878      1939 5      [RAB$C_RFA] : STATUS = RMSPOS_RFA();  
; 1879      1940 5      TES  
; 1880      1941 5        
; 1881      1942 5      END)  
; 1882      1943 4      THEN  
; 1883      1944 3      RETURN .STATUS;  
; 1884      1945 3      ! Setup record pointer (RP) to the RFA/RRV of retrieved record.  
; 1885      1946 3        
; 1886      1947 3      IRAB[IRBSW_NEXT_ID] = RMSRECORD_ID();  
; 1887      1948 3        
; 1888      1949 3      AP = 3;  
; 1889      1950 3        
; 1890      1951 3      BEGIN  
; 1891      1952 3        
; 1892      1953 4      GLOBAL REGISTER  
; 1893      1954 4      R_BDB;  
; 1894      1955 4        
; 1895      1956 4      IRAB[IRBSL_NEXT_VBN] = RMSRECORD_VBN();  
; 1896      1957 4      END:  
; 1897      1958 4      ! Move the key of the internally current record into keybuffer 2 in  
; 1898      1959 3      preparation for updating the local NRP context. It will not be necessary  
; 1899      1960 3      to extract the key, if positioning was done by means of an alternate key  
; 1900      1961 3      index, because as part of that positioning, the key would have been moved  
; 1901      1962 3      into keybuffer 2.  
; 1902      1963 3        
; 1903      1964 3      IF .IRAB[IRBSB_RP_KREF] EQLU 0  
; 1904      1965 3      THEN  
; 1905      1966 3      ! If any of the following conditions holds, the key of the internally  
; 1906      1967 3      current record (which must be a primary data record) must be  
; 1907      1968 3      extracted from the record itself.  
; 1908      1969 3        
; 1909      1970 3      1. If the file is a prologue 1 or 2 file.  
; 1910      1971 3      2. If the record was retrieved randomly.  
; 1911      1972 3      3. If key compression is not enabled in this key of reference.  
; 1912      1973 3      4. If deleted records were encountered during the positioning.  
; 1913      1974 3        
; 1914      1975 3      IF ((.IFAB[IFBSB_PLG_VER] LSSU PLG$C_VER_3)  
; 1915      1976 3      OR  
; 1916      1977 3      (.RAB[RAB$B_RAC] NEQU RAB$C_SEQ)  
; 1917      1978 3      OR  
; 1918      1979 5      NOT .IDX_DFN[DX$V_KEY_COMPR]  
; 1919      1980 4      OR  
; 1920      1981 5      .IRAB[IRBSV_DEL_SEEN])  
; 1921      1982 4      THEN  
; 1922      1983 4      BEGIN  
; 1923      1984 4        
; 1924      1985 4      GLOBAL REGISTER  
; 1925      1986 3      R_BDB;  
; 1926      1987 4        
; 1927      1988 4      AP = 0;  
; 1928      1989 4      RMSRECORD_KEY (KEYBUF_ADDR(2));  
; 1929      1990 4      END  
; 1930      1991 4        
; 1931      1992 4        
; 1932      1993 4        
; 1933      1994 4        
; 1934      1995 4      
```

```
: 1935 1996 4      ! If this is a prologue 3 file, and RMS is positioning sequentially,
: 1936 1997 4      then RMS may use the key of the last retrieved record, stored in
: 1937 1998 4      keybuffer 1 to supply any characters front compressed off the current
: 1938 1999 4      key provided no intervening records were encountered between the
: 1939 2000 4      last retrieved, and the new record.
: 1940 2001 4
: 1941 2002 3      ELSE
: 1942 2003 4          BEGIN
: 1943 2004 4
: 1944 2005 4          LOCAL
: 1945 2006 4              KEY      : REF BBLOCK;
: 1946 2007 4
: 1947 2008 4          MACRO
: 1948 2009 4              KEY_LEN = 0,0,8,0 %,
: 1949 2010 4              CMP_CNT = 1,0,8,0 %;
: 1950 2011 4
: 1951 2012 4          KEY = .REC_ADDR + RMSREC_OVHD(0);
: 1952 2013 4
: 1953 2014 4          CH$COPY (.KEY[CMP_CNT], KEYBUF_ADDR(1),
: 1954 2015 4              .KEY[KEY_LEN], .KEY+2
: 1955 2016 4              :(.KEY + .KEY[KEY_LEN] + 1),
: 1956 2017 4              .IDX_DFN[IDX$B_KEY$Z], KEYBUF_ADDR(2));
: 1957 2018 3          END:
: 1958 2019 3
: 1959 2020 3      ! Don't do any record-locking if there aren't any writers of the file, and
: 1960 2021 3      if pseudo record locking is not to be done.
: 1961 2022 3
: 1962 2023 3      IF .IFAB[IFBSV_NORECLK]
: 1963 2024 3          AND
: 1964 2025 4          NOT (.IFAB[IFBSV_RU_RLK]
: 1965 2026 4          AND
: 1966 2027 4              .IFAB[IFBSV_RUP])
: 1967 2028 3      THEN
: 1968 2029 3          RETURN .STATUS;
: 1969 2030 3
: 1970 2031 2          END;                                ! of block defining local STATUS
: 1971 2032 2
: 1972 2033 3      BEGIN
: 1973 2034 3
: 1974 2035 3      .LOCAL
: 1975 2036 3          STATUS;
: 1976 2037 3
: 1977 2038 4      BEGIN
: 1978 2039 4
: 1979 2040 4      LABEL
: 1980 2041 4          OK_WAT;
: 1981 2042 4
: 1982 2043 4      ! Flag no special action needed for unlocking the RP
: 1983 2044 4
: 1984 2045 4      AP = 0;
: 1985 2046 4
: 1986 2047 4
: 1987 2048 4      ! only query_lock the record if:
: 1988 2049 4
: 1989 2050 4      1. The user has specified no locking ( NLK ) and is not in a Recovery
: 1990 2051 4      Unit.
: 1991 2052 4
```

```
1992 2053 4 | 2. The file is opened for read ( not WRTACC )
1993 2054 4 | and the user has not specified read-only locking ( not REA )
1994 2055 4
1995 2056 5 IF (.RAB[RAB$V_NLK]
1996 2057 5 AND
1997 2058 5 NOT .IFAB[IFBS$V_RUP])
1998 2059 5
1999 2060 4 OR
2000 2061 6 ( (NOT .IFAB[IFBS$V_WRTACC])
2001 2062 5 AND
2002 2063 5 (NOT .RAB[RAB$V_REA]) )
2003 2064 4 THEN STATUS = RMSQUERY_LCK (.IRAB[IRBSL_NEXT_VBN], .IRAB[IRBSW_NEXT_ID])
2004 2065 4 ELSE STATUS = RMSLOCK (.IRAB[IRBSL_NEXT_VBN], .IRAB[IRBSW_NEXT_ID]);
2005 2066 4
2006 2067 4
2007 2068 4
2008 2069 4
2009 2070 4 | OK_WAT success status means we had to wait for someone else to unlock the
2010 2071 4 | record. To wait, we deaccessed the bucket. Therefore, we must reaccess
2011 2072 4 | it, and we can use the record pointer information for this. Deaccessing
2012 2073 4 | the bucket also means that our NRP context updating information in the
2013 2074 4 | IRAB cannot longer be considered to be valid.
2014 2075 4
2015 2076 5 IF .STATUS EQL RMSSUC(OK_WAT)
2016 2077 4 THEN
2017 2078 4 OK_WAT:
2018 2079 5 BEGIN
2019 2080 5
2020 2081 5 LOCAL
2021 2082 5 TEMP_STATUS;
2022 2083 5
2023 2084 5 | Reposition to the record using record pointer contents. If it is
2024 2085 5 | possible that some reclamation maybe done make sure the primary data
2025 2086 5 | bucket is exclusively accessed.
2026 2087 5
2027 2088 5 IF .IFAB[IFBS$V_WRTACC]
2028 2089 5 AND
2029 2090 5 .IFAB[IFBS$V_RU]
2030 2091 5 THEN
2031 2092 5 IRAB[IRBSB_CACHEFLGS] = CSHSM_LOCK;
2032 2093 5
2033 2094 6 IF NOT (TEMP_STATUS = RMSFIND_BY_RRV (.IRAB[IRBSL_NEXT_VBN],
2034 2095 6 .IRAB[IRBSW_NEXT_ID],
2035 2096 6 0))
2036 2097 5 THEN
2037 2098 5 STATUS = .TEMP_STATUS;
2038 2099 5
2039 2100 5 | If RMS after re-positioning to the record finds that it had been
2040 2101 5 deleted within a Recovery Unit, then RMS will have to re-position
2041 2102 5 after deleting the record for good if it has write access to the
2042 2103 5 file.
2043 2104 5
2044 2105 5 IF .STATUS
2045 2106 5 AND
2046 2107 5 .REC_ADDR[IRCSV_RU_DELETE]
2047 2108 5 THEN
2048 2109 5 LEAVE OK_WAT;
```

```
2049      2110 5
2050      2111 5
2051      2112 5
2052      2113 5
2053      2114 5
2054      2115 5
2055      2116 5
2056      2117 5
2057      2118 6
2058      2119 5
2059      2120 6
2060      2121 6
2061      2122 6
2062      2123 6
2063      2124 6
2064      2125 6
2065      2126 6
2066      2127 6
2067      2128 6
2068      2129 6
2069      2130 6
2070      2131 6
2071      2132 6
2072      2133 6
2073      2134 6
2074      2135 6
2075      2136 6
2076      2137 6
2077      2138 5
2078      2139 6
2079      2140 6
2080      2141 6
2081      2142 6
2082      2143 6
2083      2144 6
2084      2145 6
2085      2146 5
2086      2147 5
2087      2148 5
2088      2149 4
2089      2150 4
2090      2151 4
2091      2152 4
2092      2153 4
2093      2154 4
2094      2155 4
2095      2156 4
2096      2157 4
2097      2158 4
2098      2159 4
2099      2160 5
2100      2161 5
2101      2162 5
2102      2163 4
2103      2164 5
2104      2165 5
2105      2166 5

      IF .STATUS
      THEN
          ! If our key of reference is the primary key, then we can reclaim
          ! our NRP updating information from the primary data bucket's
          ! VBN and the record ID.
          IF (.IDX_DFNE[IDX$B_KEYREF] EQL 0)
          THEN
              BEGIN
                  IRAB[IRBSL_RFA_VBN] = .BBLOCK[.IRAB[IRBSL_CURBDB], BDBSL_VBN];
                  IRAB[IRBSW_RFA_ID] = IRCS_ID(REC_ADDR);
              END
          ! If the key of reference is not the primary key, then RMS has no
          ! easy way to reclaim the NRP list updating information which is
          ! for the SIDR bucket (long since released) and not the primary
          ! data bucket. Since the stream which has the record locked might
          ! delete the SIDR array positioned to but not the primary data
          ! record itself (by means of an $UPDATE), RMS must re-position
          ! in order to guarantee that the key of the SIDR array it
          ! positions to is actually represented in the primary data record
          ! to be returned. The alternate success status, and the fact that
          ! the key of reference is other than the primary will force another
          ! attempt to access the primary data bucket after accessing the
          ! necessary SIDR, and to lock the next record.
          ELSE
              BEGIN
                  GLOBAL REGISTER
                  R_BDB_STR;
                  RELEASE (IRAB[IRBSL_CURBDB]);
              END
          ELSE
              IRAB[IRBSL_CURBDB] = 0;
          END;
      ! If RMS finds that the current record has been modified within a Recovery
      ! Unit, then subject it to further processing before deciding whether to
      ! return it as the non-deleted primary data record to be returned to the
      ! user, or whether to return a status to force RMS to re-position.
      IF .STATUS
      AND
      .IRAB[IRBSL_CURBDB] NEQU 0
      AND
      (.REC_ADDR[IRCSV_RU_DELETE]
      OR
      .REC_ADDR[IRCSV_RU_UPDATE])
      THEN
          BEGIN
              LOCAL
```

```
2106      5
2107      5
2108      5
2109      5
2110      5
2111      5
2112      5
2113      5
2114      5
2115      5
2116      5
2117      5
2118      5
2119      5
2120      5
2121      5
2122      5
2123      5
2124      5
2125      5
2126      5
2127      5
2128      5
2129      5
2130      5
2131      5
2132      5
2133      6
2134      6
2135      6
2136      6
2137      6
2138      6
2139      7
2140      6
2141      6
2142      7
2143      6
2144      6
2145      5
2146      4
2147      4
2148      4
2149      5
2150      5
2151      5
2152      4
2153      5
2154      5
2155      5
2156      5
2157      4
2158      4
2159      4
2160      3
2161      3
2162      3

2167      5
2168      5
2169      5
2170      5
2171      5
2172      5
2173      5
2174      5
2175      5
2176      5
2177      5
2178      5
2179      5
2180      5
2181      5
2182      5
2183      5
2184      5
2185      5
2186      5
2187      5
2188      5
2189      5
2190      5
2191      5
2192      5
2193      6
2194      6
2195      6
2196      6
2197      6
2198      6
2199      6
2200      7
2201      6
2202      7
2203      6
2204      6
2205      5
2206      4
2207      4
2208      4
2209      4
2210      5
2211      5
2212      5
2213      4
2214      5
2215      5
2216      5
2217      4
2218      4
2219      4
2220      3
2221      3
2222      3
2223      3

      RECORD_ID : WORD;
      RECORD_ID = .REC_ADDR[IRC$W_ID];
      ! If the file has been open for write access, then attempt to delete
      ! the record if it was deleted within a Recovery Unit, or re-format
      ! the record if it was updated within a Recovery Unit.
      IF .IFAB[IFBSV_WRTACC]
      THEN
        RMSRU_RECLAIM();
      ! If the record had been deleted within a Recovery Unit, then RMS will
      ! not return this record to the user as a non-deleted primary data
      ! record. Therefore, release the primary data bucket, and change the
      ! return status to 0 if RMS did not have to wait for the record lock, or
      ! change the return status to RMSS_DEL if RMS had to wait for the
      ! record lock. Returning a status of RMSS_DEL in the latter case will
      ! allow the information that RMS had to wait for a record lock to
      ! eventually be returned to the user along with the non-deleted primary
      ! data record when such a record is eventually found.
      IF .RECORD_ID NEQU .REC_ADDR[IRC$W_ID]
        OR
        .REC_ADDR[IRC$V_RU_DELETE]
      THEN
        BEGIN
          GLOBAL REGISTER
          R_BDB_STR;
          RELEASE (IRAB[IRBSL_CURBDB]);
          IF .STATUS<0,16> EQLU RMSSUC(OK_WAT)
          THEN
            STATUS = RMSERR(DEL)
          ELSE
            STATUS = 0;
          END;
        END;
      ! Return here if QUERY_LCK.
      IF (.RAB[RABSV_NLK]
        AND
        NOT .IFAB[IFBSV_RUP])
        OR
        (NOT .IFAB[IFBSV_WRTACC]
        AND
        NOT .RAB[RABSV_REA])
      THEN
        RETURN .STATUS;
      END;
      ! If UNLOCK_RP was set coming here, it can only mean that this was a
      ! reaccessing of a previously automatically locked record that was not
```

```
2163 2224 3 | unlocked at the beginning of this operation to avoid a locking window.
2164 2225 3 | It will get an ok_alk status (not suc) from rm$lock. It wants to release
2165 2226 3 | the current record on potential buffer errors. The case where we don't
2166 2227 3 | want to release the now current record lock is if the status from rm$lock
2167 2228 3 | was rms$_ok_alk (i.e., not suc) which meant that it had been previously
2168 2229 3 | manually locked, and should remain that way even if this operation fails.
2169 2230 3
2170 2231 3 IF TESTBITCC([RAB[IRBSV_UNLOCK_RP])
2171 2232 3 THEN
2172 2233 4 | IF .STATUS<0,16> EQLU RMSSUC()
2173 2234 3 | OR
2174 2235 4 | .STATUS<0,16> EQLU RMSSUC(OK_WAT)
2175 2236 3 | OR
2176 2237 4 | .STATUS<0,16> EQLU RMSSUC(OK_RULK)
2177 2238 3 THEN
2178 2239 3 | AP = 1
2179 2240 3 ELSE
2180 2241 3
2181 2242 3 | If it is necessary for us to release the record lock set
2182 2243 3 | IRBSV_UNLOCK_RP. It will only be necessary to release the record
2183 2244 3 | lock in some circumstances when we have had to stall waiting for
2184 2245 3 | it, and whenever RMS has managed to position to a record that was
2185 2246 3 | deleted within a Recovery Unit. In the former case if after
2186 2247 3 | waiting for the lock, we have some problem reaccessing the bucket,
2187 2248 3 | or we find that the record is deleted while we were waiting, then
2188 2249 3 | we must release the record lock. If we are positioning by means
2189 2250 3 | of an alternate key and we have had to stall waiting for the
2190 2251 3 | record lock, and this is an operation where the NRP list must be
2191 2252 3 | updated (any operation but a nonrandom SFIND), then the record
2192 2253 3 | lock must also be released.
2193 2254 3
2194 2255 4 BEGIN
2195 2256 4
2196 2257 5 IF (.IRAB[IRBSL_CURBDB] EQL 0)
2197 2258 4 | AND
2198 2259 5 | (.STATUS<0,16> NEQ RMSERR(RLK))
2199 2260 4 THEN
2200 2261 4 | IRAB[IRBSV_UNLOCK_RP] = 1;
2201 2262 4
2202 2263 4 RETURN .STATUS;
2203 2264 4 END
2204 2265 3 ELSE
2205 2266 3 | AP = 1;
2206 2267 3
2207 2268 3 IF NOT .RAB[RAB$V_ULK]
2208 2269 3 THEN
2209 2270 3 | IRAB[IRBSV_UNLOCK_RP] = 1;
2210 2271 3
2211 2272 3 RETURN .STATUS;
2212 2273 3
2213 2274 3 END
2214 2275 1 END: | of local block defining STATUS
2275 1 | of routine
```

		5E	18	C2 00000 GET_RECORD:					
		40	A9	94 00003	SUBL2	#24, SP		1427	
		42	A9	B4 00006	CLRB	64(IRAB)		1507	
		04	A9	9E 00009	CLRW	66(IRAB)		1508	
		14	AE	1000	8F AA 0000E	MOVAB	4(IRAB), 20(SP)	1509	
		14	BE	38	A8 D4 00014	BICW2	#4096, 20(SP)		
		02	00	1E	A8 8F 00017	CLRL	56(RAB)	1515	
	01A3	0097		0000D	0001C 1\$:	CASEB	30(RAB), #0, #2	1524	
					.WORD	2\$-1\$,-			
						12\$-1\$,-			
						29\$-1\$			
		50	8644	8F 3C 00022	MOVZWL	#34372, R0		1862	
				4C 11 00027	BRB	8\$			
		42	14	BE	17 E0 00029	2\$::	BBS	#23, 20(SP), 7\$	
		38	14	BE	05 E1 0002E	BBC	#5, 20(SP), 6\$	1549	
		33	14	BE	09 E0 00033	BBS	#9, 20(SP), 6\$	1553	
		0C	14	BE	0B E5 00038	BBCC	#11, 20(SP), 3\$	1555	
		2E	14	BE	01 E0 00030	BBS	#1, 20(SP), 7\$	1563	
				C9 9A 00042	MOVZBL	195(IRAB), -(SP)	1570		
				08 11 00047	BRB	4\$	1574		
				FC87 30 00049	RSBW	SETUP_NRP_DATA			
				0000G 30 00051	MOVZBL	194(IRAB), -(SP)	1584		
				04 C0 00054	BSBW	RMSKEY_DESC	1585		
		06	00A6	C9	20 A7 90 00057	ADDL2	#4, SP		
		06	00A2	A8	04 E1 00050	MOVB	32(IDX_DFN), 166(IRAB)	1588	
				02 CA	02 E1 00062	BBC	#4, 6(RAB), 5\$	1596	
		08	14	BE	01A0 31 00068	BRW	#2, 162(IFAB), 11\$	1598	
				50	55: 31 0006B	BBC	36\$	1600	
				827A	65: 30 00070	MOVZWL	#33402, R0	1608	
				03B6	75: 31 00075	BRW	71\$	1610	
				0800	85: A8 00078	BISW2	#2048, 20(SP)		
				00A8	95: D5 0007E	TSTL	168(IRAB)		
					06 12 00082	BNEQ	10\$		
				0800	AA 00084	BICW2	#2048, 20(SP)	1620	
				7E	00C3 0000G	MOVZBL	195(IRAB), -(SP)	1622	
					30 0008F	BSBW	RMSKEY_DESC		
					04 C0 00092	ADDL2	#4, SP		
				00A6	5E DD	50 E9 00095	BLBC	STATUS, 8\$	
					20 A7 90 00098	MOVB	32(IDX_DFN), 166(IRAB)	1624	
					20 A7 9A 0009E	MOVZBL	32(IDX_DFN), R1	1625	
					00B4 50 CA	3C 000A2	MOVZWL	180(IFAB), R0	
					60 A9	C0 000A7	ADDL2	96(IRAB), R0	
					51 28 000AB	MOVC3	R1, 296(IRAB), (R0)		
					0135 31 000B0	BRW	33\$		
					11\$: 00 F0 000B3	INSV	#0, #23, #1, 20(SP)	1524	
					12\$: 8F AA 000B9	BICW2	#2048, 20(SP)	1638	
					35 A8 9A 000BF	MOVZBL	53(RAB), -(SP)	1640	
					0000G 30 000C3	BSBW	RMSKEY_DESC	1642	
					04 C0 000C6	ADDL2	#4, SP		
					50 E9 000C9	BLBC	STATUS, 8\$		
					34 A8 90 000CC	MOVB	52(RAB), KEYSIZE	1643	
					1D A7 95 000D0	TSTB	29(IDX_DFN)	1647	
					06 13 000D3	BEQL	13\$		
					01 1E A7 91 000D5	CMPB	30(IDX_DFN), #1	1648	
					0C 1B 000D9	BLEQU	14\$		
					50 95 000DB	TSTB	KEYSIZE	1652	

20	A7	16	13 000DD	BEQL	16\$		1654
		50	91 000DF	CMPB	KEYSIZE, 32(IDX_DFN)		
		17	18 000E3	BLEQU	17\$		1656
		0E	11 000E5	BRB	16\$		1662
		50	95 000E7	TSTB	KEYSIZE		
		04	12 000E9	BNEQ	15\$		
20	50	20	A7 90 000EB	MOV8	32(IDX_DFN), KEYSIZE		1664
			50 91 000EF	CMPB	KEYSIZE, 32(IDX_DFN)		1666
	50	85A4	07 13 000F3	BEQL	17\$		1668
			4F 11 000FA	BRB	20\$		
62	00A6	C9	50 90 000FC	MOV8	KEYSIZE, 166(IRAB)		1672
		52	30 A8 D0 00101	MOVL	48(RAB), KBF_ADDR		1673
		50	0A A9 0C 00105	PROBER	10(IRAB), KEYSIZE, (KBF_ADDR)		1675
			07 12 0010A	BNEQ	18\$		
		50	858C	MOVZWL	#34212, R0		
			8F 3C 0/10C	BRB	20\$		
60			2D 11 C0111	MOVZWL	#34188, R0		1679
		51	50 9A 00113	MOVZBL	KEYSIZE, R1		
		50	00B4 CA 3/ 00116	MOVZWL	180(IFAB), R0		
		50	60 A9 00 0011B	ADDL2	96(IRAB), R0		
	62	05	51 28 0011F	MOVC3	R1, (KBF_ADDR) (R0)		1683
		05	1D A7 91 00123	CMPB	29(IDX_DFN), #5 (R0)		
			06 12 00127	BNEQ	19\$		
			00000G 30 00129	BSBW	RMS\$PCKDEC CHECK		1685
12		11	50 E9 0012C	BLBC	STATUS, 20\$		
08		54	04 A8 9E 0012F	MOVAB	4(RAB), R4		1689
		64	15 E1 00133	BBC	#21, (R4), 22\$		
		64	16 E1 00137	BBC	#22, (R4), 21\$		1692
		50	867C 8F 3C 00138	MOVZWL	#34428, R0		1694
			02EB 31 00140	BRW	71\$		
07		42	A9 10 88 00143	BISB2	#16, 66(IRAB)		1697
			08 11 00147	BRB	23\$		1698
	42	64	16 E1 00149	BBC	#22, (R4), 24\$		1703
		A9	02 88 0014D	BISB2	#2 66(IRAB)		1706
F8		04	A9 0094 31 00151	BRW	33\$		1707
			05 E0 00154	BBS	#5, 4(IRAB), 23\$		1729
			21 A7 95 00159	TSTB	33(IDX_DFN)		1731
			F3 12 0015C	BNEQ	23\$		
		20	EF A7 00A6 C9 91 00162	BLBS	28(IDX_DFN), 23\$		1733
			7E 12 J0168	CMPB	166(IRAB), 32(IDX_DFN)		1737
			51 00A6 C9 9A 0016A	MOVZBL	166(IRAB), SIZE		
			50 00B4 CA 3C 0016F	MOVZWL	180(IFAB), R0		1746
			60 B940 3F 00174	PUSHAW	#96(IRAB)[R0]		1748
9E		60 B940	51 29 00178	CMPC3	SIZE, #96(IRAB)[R0], a(SP)+		
			68 12 0017E	BNFQ	33\$		
		51	00B0 C9 D0 00180	MOVL	176(IRAB), VBN		1758
			61 13 00185	BEQL	33\$		
0C		06	AA 03 E1 00187	BRC	#3, 6(IFAB), 25\$		1769
27	00A2	CA 03 E1 0018C	BBC	#3, 162(IFAB), 28\$		1771	
21	00A2	CA 02 E1 00192	BBC	#2, 162(IFAB), 28\$		1773	
06		64 14 E1 00198	25\$:	BBC	#20, (R4), 26\$		1780
46	00A2	CA 02 E1 0019C	BBC	#2, 162(IFAB), 33\$		1782	
04		64 11 E1 001A2	26\$:	BBC	#17, (R4), 27\$		1793
	07	A9 01 88 001A6	BISB2	#1, 7(IRAB)			1795
		52 00BC C9 3C 001AA	MOVZWL	188(IRAB), R2			1797
		0000C 30 001AF	BSBW	RMS\$QUERY_LCK			

14	BE	01	8039	8F	50	81	001B2	CMPW	ST	#32825	1799	
			05	A9	2F	12	001B7	BNEQ	33\$		4	
			14	BE	10	88	001B9	28\$:	BISB2	#16, 5(IRAB)	1835	
					4C	11	001BD	BRB	36\$		4	
			14	17	00	F0	001BF	29\$:	INSV	#0, #23, #1, #20(SP)	1836	
				14	BE	8F	AA	001C5	BICW2	#2048, #20(SP)	1847	
					10	A8	05	001CB	TSTL	16(RAB)	1849	
						05	13	001CE	BEQL	30\$	1851	
					14	A8	B5	001D0	TSTW	20(RAB)	1853	
						08	12	001D3	BNEQ	32\$	1855	
			50		8F	3C	001D5	30\$:	MOVZWL	#34396, R0	1857	
					0251	31	001DA	31\$:	BRW	71\$	1857	
					7E	D4	001DD	32\$:	CLRL	-(SP)	1857	
					0000G	30	001DF	BSBW	RMSKEY_DESC		1857	
					04	C0	001E2	ADDL2	#4, SP		1857	
			1E	04	5E	50	E9	001E5	BLBC	STATUS, 31\$	1871	
				08	F2	0D	E5	001E8	BBCC	#13, 4(IRAB), 36\$	1879	
				52	0080	1C	AE	E9	BLBC	REPOS STATUS, 34\$	1881	
				51	78	C9	3C	001F1	MOVZWL	128(IRAB), R2		
						A9	D0	001F6	MOVL	120(IRAB), R1		
						0C	11	001FA	BRB	35\$		
					51	00B0	C9	D0	001FC	MOVL	176(IRAB), R1	1883
						08	13	00201	BEQL	36\$		
					52	00BC	C9	3C	00203	MOVZWL	188(IRAB), R2	1885
						0000G	30	00208	BSBW	RMSUNLOCK		
			0A	05	00C2	21	A7	90	0020B	MOV8	33(IDX DFN), 194(IRAB)	1888
				50	50	1E	A8	90	00211	MOVB	30(RAB), RAC	1904
				50	00C3	04	E1	00215	BBC	#4, 5(IRAB), 37\$	1906	
						C9	90	0021A	MOVB	195(IRAB), RAC	1918	
						03	13	0021F	BEQL	37\$	1920	
			02	00	50	02	90	00221	MOVB	#2, RAC	1922	
			0010	0008	50	50	8F	00224	CASEB	RAC, #0, #2	1926	
						0006	00228	37\$:	.WORD	39\$-38\$,-		
										40\$-38\$,-		
										41\$-38\$		
						0000G	30	0022E	39\$:	BSBW	RMSPOS_SEQ	1931
						08	11	00231	BRB	42\$		
						0000G	30	00233	40\$:	BSBW	RMSPOS_KEY	1935
						03	11	00236	BRB	42\$		
						0000G	30	00238	41\$:	BSBW	RMSPOS_RFA	1939
			14	AE	14	50	D0	0023B	42\$:	MOVL	R0, STATUS	
				03		AE	E8	0023F		BLBS	STATUS, 43\$	1926
						00AE	31	00243	BRW	48\$		
						0000G	30	00246	43\$:	BSBW	RMSRECORD_ID	1949
			0080	C9		50	B0	00249	MOVW	R0, 128(IRAB)		
				5C		03	D0	0024E	MOVL	#3, AP	1951	
						0000G	30	00251	BSBW	RMSRECORD_VBN	1958	
			78	A9		50	D0	00254	MOVL	R0, 120(IRAB)		
						00C2	C9	95	TSTB	194(IRAB)		
						27	12	00258	BNEQ	45\$		
						03	0F	0025C	CMPB	183(IFAB), #3	1979	
						03	CA	91	PLSSU	44\$		
						1E	A8	95	TSTB	30(RAB)	1981	
						06	E1	00263	BNEQ	44\$		
						01	E1	0026A	BBC	#6, 28(IDX DFN), 44\$	1983	
						5C	D4	00274	CLRL	#1, 67(IRAB), 46\$	1985	
									AP		1992	

3C	66	05	E0 0035E	BBS	#5, (REC_ADDR), 60\$: 2107	
	36	52	E9 00362	BLBC	STATUS, 59\$: 2121	
		21	A7 95 00365	TSTB	33(IDX_DFN)	: 2118	
		20	12 00368	BNEQ	58\$		
70	50	20	A9 D0 0036A	MOVL	32(IRAB), R0	: 2121	
	A9	1C	A0 D0 0036E	MOVL	28(R0), 112(IRAB)		
	03	00B7	CA 91 00373	CMPB	183(IFAB), #3	: 2122	
		06	1E 00378	BGEQU	56\$		
	50	01	A6 9A 0037A	MOVZBL	1(REC_ADDR), R0		
		04	11 0037E	BRB	57\$		
74	50	01	A6 3C 00380	MOVZWL	1(REC_ADDR), R0		
	A9	50	B0 00384	MOVW	R0, 1T6(IRAB)		
		14	11 00388	BRB	60\$: 2118	
	54	20	A9 D0 0038A	MOVL	32(IRAB), BDB	: 2144	
		20	A9 D4 0038E	CLRL	32(IRAB)		
		7E	D4 00391	CLRL	-(SP)		
		0000G	30 00393	BSBW	RMSRLSBKT		
	5E	04	C0 00396	ADDL2	#4, SP		
		03	11 00399	BRB	60\$: 2118	
	44	20	A9 D4 0039B	CLRL	32(IRAB)	: 2147	
		52	E9 0039E	BLBC	STATUS, 65\$: 2156	
		20	A9 D5 003A1	TSTL	32(IRAB)	: 2158	
		3F	13 003A4	BEQL	65\$		
04	66	05	E0 003A6	BBS	#5, (REC_ADDR), 61\$: 2160	
	66	06	E1 003AA	BBC	#6, (REC_ADDR), 65\$: 2162	
	53	01	A6 B0 003AE	MOVW	1(REC_ADDR), RECORD_ID	: 2169	
	06	06	AA E9 003B2	BLBC	6(IFAB), 62\$: 2175	
	01	A6	00000000G	EF 16 003B6	JSB	RMSRU RECLAIM	: 2177
		53	B1 003BC	CMPW	RECORD_ID, 1(REC_ADDR)	: 2189	
	04	12 003C0	BNEQ	63\$			
1F	66	05	E1 003C2	BBC	#5, (REC_ADDR), 65\$: 2191	
	54	20	A9 D0 003C6	MOVL	32(IRAB), BDB	: 2198	
		20	A9 D4 003CA	CLRL	32(IRAB)		
	7E	D4 003CD	CLRL	-(SP)			
	0000G	30 003CF	BSBW	RMSRLSBKT			
	5E	04	C0 003D2	ADDL2	#4, SP		
8061	8F	52	B1 003D5	CMPW	STATUS, #32865	: 2200	
		07	12 003DA	BNEQ	64\$		
	52	8262	8F 3C 003DC	MOVZWL	#33378, STATUS	: 2202	
		02	11 003E1	BRB	65\$		
		52	D4 003E3	CLRL	STATUS	: 2204	
06	06	A8	04 E1 003E5	BBC	#4, 6(RAB), 66\$: 2210	
3B	00A2	CA	02 E1 003EA	BBC	#2, 162(IFAB), 70\$: 2212	
	05	06	AA E8 003F0	BLBS	6(IFAB), 67\$: 2214	
32	04	A8	02 E1 003F4	BBC	#2, 4(RAB), 70\$: 2216	
21	04	A9	0D E4 003F9	BBSC	#13, 4(IRAB), 68\$: 2231	
	01		52 B1 003FE	CMPW	STATUS, #1	: 2233	
		1C 13 00401	BEQL	68\$			
8061	8F	52	B1 00403	CMPW	STATUS, #32865	: 2235	
		15 13 00408	BFQL	68\$			
8071	8F	52	B1 0040A	CMPW	STATUS, #32881	: 2237	
		0E 13 0040F	BEUL	68\$			
		20 A9 D5 00411	TSTL	32(IRAB)	: 2257		
82AA	8F	52	B1 00416	BNEQ	70\$: 2259	
		0A 12 00418	CMPW	STATUS, #33450			
		0C 11 0041D	BNEQ	69\$			
			BPB	70\$: 2263		

04	06	5C	01	D0 0041F 68\$:	MOVL #1, AP	2266
	05	A8	02	E0 00422	BBS #2, 6(RAB)	2268
		A9	20	88 00427 69\$:	BISB2 #3, 5(IRA8)	2270
		50	52	D0 0042B 70\$:	MOVL STATUS, R0	2272
		5E	18	C0 0042E 71\$:	ADDL2 #24, SP	2275
				05 00431	RSB	:

; Routine Size: 1074 bytes, Routine Base: RMS\$RMS3 + 032D

: 2215 2276 1 END ! of module
: 2216 2277 1
: 2217 2278 0 ELUDOM

PSECT SUMMARY

Name	Bytes	Attributes
RMS\$RMS3	1887	NOVEC,NOWRT, RD, EXE,NOSHR, GBL, REL, CON, PIC,ALIGN(2)

Library Statistics

File	----- Symbols -----			Pages Mapped	Processing Time
	Total	Loaded	Percent		
\$_\$255\$DUA28:[RMS.OBJ]RMS.L32;1	3109	140	4	154	00:00.4

COMMAND QUALIFIERS

: BLISS/CHECK=(FIELD,INITIAL,OPTIMIZE)/LIS=LIS\$:RM3GET/OBJ=OBJ\$:RM3GET MSRC\$:RM3GET/UPDATE=(ENH\$:RM3GET)

: Size: 1887 code + 0 data bytes
: Run Time: 00:48.3
: Elapsed Time: 01:24.0
: Lines/CPU Min: 2829
: Lexemes/CPU-Min: 17134
: Memory Used: 399 pages
: Compilation Complete

0325 AH-BT13A-SE
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